

The Berlekamp-Hensel Procedure  
(Hans Zassenhaus ~ 1972)

Input  $A \in \mathbb{Z}[x]$  s.t.  
 $d > 1$

$$A = a_d x^d + \dots + a_0$$

$$\text{cont}(A) = 1$$

$$\text{gcd}(A, A') = 1$$

Output  $f_1, f_2, \dots, f_m \in \mathbb{Z}[x]$  s.t.

$f_i$  irreducible over  $\mathbb{Q}$

$$A = f_1 f_2 \dots f_m$$

↑  
gives  $\text{lc}(A)$ .

④ Test if  $\text{pp}(\text{ad} \cdot g_i^{(n)} \bmod p^n) \mid A \ \forall i$   
Test if  $\text{pp}(\text{ad} \cdot g_i^{(n)} \cdot g_j^{(n)} \bmod p^n) \mid A \ \forall i \neq j$   
etc.

↑  
expand

① Pick  $p$  s.t.

$$p \nmid \text{lc} A = a_d$$

$$\text{gcd}(A, A') = 1 \text{ in } \mathbb{Z}_p[x]$$

$\mathbb{F}_p$

③ For  $j = 1, 2, \dots, l$  Hensel lift  $g_j$

using  $u_0 = g_j, w_0 = \text{ad} \cdot \prod_{i \neq j} g_i$  until

$p^n > 2 \text{ad} \cdot \|f_i\|_\infty$  to obtain

$$A \equiv \text{ad} \cdot \underbrace{g_1^{(n)}}_{\text{monic}} \cdot \underbrace{g_2^{(n)}} \cdot \dots \cdot \underbrace{g_l^{(n)}} \pmod{p^n}$$

↑  
monic.

② Factor  $A \in \mathbb{Z}_p[x]$

$$A \in \mathbb{Z}_p[x] \xrightarrow{\text{Cantor-Zassenhaus}} A \equiv \text{ad} \cdot g_1 g_2 \dots g_l \pmod{p^n}$$

$O(d^3 \log^3 p)$

$$A \equiv \text{ad} \cdot g_1 g_2 \dots g_l \pmod{p^n} \quad l \geq m$$

$g_i \in \mathbb{Z}_p[x], \text{monic, irreducible}$

① Choose one large prime  $p > 2 \text{ad} \cdot \|f\|_\infty$  to avoid H.L.?  
This increases the cost of step ②.

② Choose many small primes  $p_i$  and apply CRT.?  
This results in too many combinations.

③ Factor mod  $\sim \ln \deg a$  primes and use the factorization  
with the least # of factors for the H.L.s

④ In 2002 Mark van Hoeij solved the combinatorial search  
for factors in step ④.

Factor the following polynomial over  $\mathbb{Z}$  by first factoring it modulo a suitably chosen prime  $p$  and employing linear Hensel lifting.

```
> a := 35*x^4+77*x^2+51*x-15*x^3-36;
      a := 35 x^4 - 15 x^3 + 77 x^2 + 51 x - 36
```

```
> gcd(a,diff(a,x));
```

1

```
> content(a,x);
```

1

```
> `mod` := mods;
```

mod := mods

```
> Factor(a) mod 2;
```

$x(x+1)^3$

← NOT SQUAREFREE

```
> Factor(a) mod 3;
```

$(-x^2)(x^2+1)$

```
> Factor(a) mod 11;
```

$2(x+4)(x^2-4x+5)(x-2)$

```
> Factor(a) mod 13;
```

$-4(x+3)(x^2-3x+6)(x-6)$

We cannot use  $p=2$  nor  $p=3$  since the polynomial is not square-free modulo those primes.

Let us use  $p=11$  noting that the factorization modulo 11 is square-free hence is assured.

```
> p := 11;
```

$p := 11$

```
> alpha := lcoeff(a,x);
```

$\alpha := 35$

The Mignotte bound on the the biggest coefficient of any factor of  $a(x)$

```
> d := degree(a,x);
```

```
  B := ceil( alpha*maxnorm(a)*2^degree(a,x)*sqrt(d+1) );
```

$B := 96420$

```
> p^4-B, p^5-B;
```

$-81779, 64631$

```
> DiophantSolve := proc(a,b,c,x,p)
```

```
  local g,sigma,tau,q,s,t;
```

```
  g := Gcdex(a,b,x,'s','t') mod p;
```

```
  if g <> 1 then error "a and b are not relatively prime!" fi;
```

```
  sigma := Rem(c*s,b,x,'q') mod p;
```

```
  # c s a = b (aq) + sigma a
```

```
  tau := Expand(c*t+q*a) mod p;
```

```
  return( sigma,tau );
```

```
end;
```

Let us lift the first factor  $x - 2$  up to the bound

```
> u[0] := x-2 mod p; ← Lift  $u_0$ 
```

$$u_0 := x - 2$$

```
> w[0] := Expand( alpha*(x+4)*(x^2-4*x+5) ) mod 11;
```

$$w_0 := 2x^3 - 4$$

```
> U := u[0];
W := w[0];
for k while p^k < 2*B do
  e[k] := expand( a-U*W );
  if k=1 then print(evaln(e[k])=e[k]); fi;
  if e[k]=0 then break; fi;
  c[k] := (e[k]/p^k) mod p;
  u[k], w[k] := DiophantSolve( w[0], u[0], c[k], x, p );
  U := U + u[k]*p^k;
  W := W + w[k]*p^k;
od;
```

$$U := x - 2$$

$$W := 2x^3 - 4$$

$$e_1 = 33x^4 - 11x^3 + 77x^2 + 55x - 44$$

```
> 'U' = U, 'W' = W;
```

$$U = x + 759240, W = 35x^3 + 77x + 84$$

Check that we have  $a - U \cdot W = 0 \pmod{p^k}$ .

```
> Expand( a - U*W ) mod p^k;
```

0

In principal we would lift the other factors but perhaps we have a real factor already.

Notice U is monic and  $\text{lc}(W) = \alpha = 35$ .

```
> f := alpha*U mod p^k;
```

$$f := 35x - 15$$

*cont(f)=5*

```
> f := primpart(f);
```

$$f := 7x - 3$$

```
> divide(a, f, 'g');
```

true

Thus we have found the factorization

```
> a = f*g;
```

$$35x^4 - 15x^3 + 77x^2 + 51x - 36 = (7x - 3)(5x^3 + 11x + 12)$$

But we do not know that g is irreducible because it has a non-trivial factorization modulo p.

```
> Factor(g) mod p;
```

$$\overbrace{5(x^2 - 4x + 5)}^{w_0} \overbrace{(x+4)}^{u_0}$$

Let us lift  $x + 4$  the other linear factor with  $a/(x+4)$ .

```
> u[0] := x+4 mod p;
```

$$u_0 := x + 4$$

```
> w[0] := Quo( a, u[0], x ) mod p;
```

$$w_0 := 2x^3 - x^2 + 4x + 2$$

```
> U := u[0];
```

```
W := w[0];
```

```
for k while p^k < 2*B do
```

```
  e[k] := expand( a-U*W );
```

```
  if e[k]=0 then break; fi;
```

```
  c[k] := (e[k]/p^k) mod p;
```

```
  u[k], w[k] := DiophantSolve( w[0], u[0], c[k], x, p );
```

```
  U := U + u[k]*p^k;
```

```
  W := W + w[k]*p^k;
```

```
od;
```

```
'U'=U; 'W'=W;
```

$$U := x + 4$$

$$W := 2x^3 - x^2 + 4x + 2$$

$$U = x - 159661$$

$$W = 35x^3 + 273437x^2 + 647211x - 797608$$

```
> h := alpha*U mod p^k;
```

$$h := 35x - 273452$$

```
> h := primpart(h);
```

$$h := 35x - 273452$$

```
> divide(a,h);
```

*false*

Therefore since there are no linear factors dividing the factor  $g$  determined earlier  $g$  must be irreducible (over  $\mathbb{Z}$ ) hence

```
> a = f*g;
```

$$35x^4 - 15x^3 + 77x^2 + 51x - 36 = (7x - 3)(5x^3 + 11x + 12)$$

```
> factor(a);
```

$$(7x - 3)(5x^3 + 11x + 12)$$