MATH 895, Assignment 3, Summer 2013

Instructor: Michael Monagan

Please hand in the assignment by 9:30am on June 20th.
Late Penalty -20% off for up to 24 hours late, zero after that.

Question 1: The Bareiss/Edmonds Algorithm

Reference: Ch. 9 of *Algorithms for Computer Algebra* by Geddes, Czapor and Labahn.

Part (a) (30 marks)

For an n by n matrix A with integer entries, implement ordinary Gaussian elimination and the Bareiss/Edmonds algorithms as the Maple procedures `GaussElim(A,n,'B');` and `Bareiss(A,n,'B');` to compute $\det(A)$. The algorithms should initially assign `B` a copy of the matrix A so that after the algorithm finishes and returns $\det(A)$ the value of B will be $A^{(n-1)}$. Note, you will need to take care of pivoting: if at any step k, the matrix entry $B_{k,k} = 0$ and $B_{i,k} \neq 0$ for some $k < i \leq n$, interchange row k with row i before proceeding. And remember interchanging two rows of a matrix changes the sign of the determinant.

Use `iquo(a,b)` to compute the quotient of a divided by b. When you are debugging, print out the matrices $A^{(1)}$, $A^{(2)}$, ... after each step of the elimination.

Execute both algorithms on the following matrices for $n = 2, 3, 4, ..., 10$.

```
> n := 4;
                                  n := 4
> m := 4:
> c := rand(10^m):
> A := Matrix(n,n,c);
```

$$A := \begin{bmatrix} 7926 & 8057 & 5 & 3002 \\ 2347 & 9765 & 3354 & 5860 \\ 6906 & 5281 & 5393 & 1203 \\ 311 & 9386 & 9810 & 5144 \end{bmatrix}$$

For $n = 4$ print out final triangular matrix for both algorithms.
Finally, in class we showed that $|\det(A)| < \sqrt{n}^n B^{mn}$ where $B = 10$ and $m = 4$ here.
Check this for $n = 2, 3, 4, ..., 10$.

Part (b) (10 marks)

Let F be a field, $D = F[x]$ and A be an n by n matrix over D. If we assume $\deg A_{i,j} \leq d$ and classical quadratic algorithms are used for polynomial multiplication and exact division in $F[x]$, how many arithmetic operations in F does the Bareiss/Edmonds algorithm do?

Try to get an exact formula in terms of n and d assuming $\deg A_{i,j} = d$. I suggest you do this for a 3x3 matrix first. Recall that to divide a polynomial in $F[x]$ of degree d by a polynomial of degree $m \leq d$, the classical division algorithm does at most $(d - m + 1)m$ multiplications in F.

Question 2: Solving $Ax = b$ using p-adic lifting.

Part (a) (20 marks)

Let $A \in \mathbb{Z}^{n \times n}$ and $b \in \mathbb{Z}^n$. In class I presented an algorithm for solving $Ax = b$ for $x \in \mathbb{Q}^n$ using linear $p-$adic lifting and rational number reconstruction. Implement the algorithm in Maple as the procedure `PadicLinearSolve(A,b)`. Your procedure should return the solution vector x and also print out the number of lifting steps k that are required. Test your implementation on the following examples. The first has large rationals in the solution vector. The second is constructed so that the solution vector x has very small rationals.

```
> with(LinearAlgebra):
> B := 2^16; m := 3; U := rand(B^m);
> n := 50;
> A := RandomMatrix(n,n,generator=U);
> b := RandomVector(n,generator=U);
> x := padicLinearSolve(A,b);
> convert( A.x-b, set ); # should be {0}
> y := [1,0,-1/2,2/3,4,3/4,-2,-3,0,-1];
> x := Vector( [seq( op(y), i=1..5 )] );
> b := A.x;
> b := 12*b; A := 12*A; # clear fractions
> x := padicLinearSolve(A,b);
> convert( A.x-b, set ); # should be {0}
```

To compute $A^{-1} \bmod p$ use the Maple command `Inverse(A) mod p`.
To multiply A times a vector x use `A.x` in Maple.
For rational number reconstruction use the Maple command `iratrecon`.

Part (b) (10 marks)

Suppose $\dim A = n$, $\dim b = n$ and $|A_{i,j}| < B^m$ and $|b_i| < B^m$, i.e., the coefficients in the linear system are m base B digits (or less). Suppose the $p - adic$ lifting algorithm does L

lifting steps, i.e. solves $Ax = b \bmod p^L$ and then successfully reconstructs $x \in \mathbb{Q}^n$ using rational reconstruction.

What is the running time of the algorithm assuming classical algorithms are used for integer arithmetic, rational reconstruction and matrix inverse. Express your answer in the form $O(f(m, n, L))$.

Since the integers in the solution vector x may be as large as mn base B digits, as illustrated by the first example, $L \in O(mn)$ in general. What is the running time for $L \in O(mn)$? Recall that we showed in class that the modular algorithm cost $O(mn^4 + m^2n^3)$ to solve $Ax = b$.