

# Algorithms for Additive and Projective Polynomials

Mark Giesbrecht

with

Joachim von zur Gathen and Konstantin Ziegler



Symbolic Computation Group  
Cheriton School of Computer Science  
University of Waterloo  
Waterloo, Ontario, Canada



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# Polynomial Composition and Decomposition

## Functional Composition

Let  $g, h \in F[x]$ , for a field  $F$ .

*Compose*  $g, h$  as functions  $f(x) = g(h(x)) = g \circ h$

A (generally) non-distributive operation (but not always):

$$g(h_1(x) + h_2(x)) \neq g(h_1(x)) + g(h_2(x))$$

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## Decomposition

Given  $f \in F[x]$ , can it be decomposed?

Do there exist  $g, h \in F[x]$  such that  $f = g \circ h$ ?

$$f = x^4 - 2x^3 + 8x^2 - 7x + 5$$

$$g = x^2 + 3x - 5 \quad h = x^2 - x - 2$$

$$\Rightarrow f = g \circ h$$

## Tame and Wild Decomposition

Let  $F$  be a field of characteristic  $p$  and  $f \in F[x]$  monic of degree  $n$ .

Normalize  $f, g, h$  to monic and *original*:  $h(0) = 0$

- $f$  is *tame* if  $p \nmid n$
- $f$  is *wild* if  $p \mid n$

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### Tame decomposition (mathematically)

- Ritt (1922) describes all tame decompositions and “ambiguities”.
- For a fixed  $s$ , there are either 0 or 1 monic  $h \in F[x]$  of degree  $s$  with  $h(0) = 0$  such that  $f(x) = g(h(x))$ .

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### Wild decomposition (mathematically)

- Life is much more difficult
- (Giesbrecht, 1988) For a finite field  $F$  of characteristic  $p$ , there are  $f \in F[x]$  of degree  $n$  with  $> n^{\lambda \log n}$  monic, original,  $h \in F[x]$  of degree  $s \approx \sqrt{n}$  such that  $f(x) = g(h(x))$ , where  $\lambda = (6 \log p)^{-1}$ .

## Tame and Wild Decomposition

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### Wild decomposition (mathematically)

- On the bright side, there are at most  $(n-1)/(s-1)$  *indecomposable* monic, original  $h \in F[x]$  of degree  $s$  such that  $f(x) = g(h(x))$ .  
(Von zur Gathen, Giesbrecht, Ziegler, 2010)

## Algorithms for Polynomial Decomposition

Barton & Zippel (1982)

Based on factorization of bivariate polynomials

$$f = g \circ h \iff h(x) - h(y) \mid f(x) - f(y)$$

Works as long as you can factor. Potentially exponential time.



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### Kozen & Landau (1987)

First polynomial-time algorithm for *tame* case. Noticed that the high-order coefficients of  $f$  do not depend on (monic)  $g$ .

➡ find  $h$ , then  $g$ .

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### Giesbrecht & May (2004)

Except for a very special case (Dickson polynomials), easily handled, Barton & Zippel's algorithm *runs in polynomial time!*

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## Theorem: Fried (1970) – Schur's Conjecture

Let  $f \in \mathbb{Q}[x]$  be indecomposable of degree  $n > 1$ .

- If  $n$  is not an odd prime, then  $(f(x) - f(y))/(x - y)$  is absolutely irreducible;
- If  $n$  is an odd prime, and it is not the case that  $f(x) = \alpha D_n(a, x + b) + \beta$  for  $\alpha, \beta, a, b \in \mathbb{Q}$ , where  $a = 0$  if  $n = 3$ , then  $(f(x) - f(y))/(x - y)$  is absolutely irreducible.

**Indecomposability** ➔ **Dickson or Irreducible (G & May 2005)**

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Von zur Gathen (1988,1990)

Nearly linear time decomposition in tame case.

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### Zippel (1991): Polynomial decomposition via Galois theory

If  $f = g \circ h$  then there exists a field  $L$  such that

$$\mathbb{F}(f(x)) \subsetneq L \subsetneq \mathbb{F}(x),$$

and  $L = \mathbb{F}(h(x))$  for some  $h \in \mathbb{F}[x]$ .

Find subfields by adapting Landau & Miller's (1985) algorithm to find subfields between  $\mathbb{Q}$  and  $\mathbb{Q}(\alpha)$  (for algebraic  $\alpha$ ).

➔ Polynomial time, at least in principle

## Additive Polynomials

Additive or linearized polynomials are those such that

$$f(x + y) = f(x) + f(y)$$

Non-linear additive polynomials only exist in  $F[x]$  if  $F$  has prime characteristic  $p$ , and have the form

$$f = a_0x + a_1x^p + a_2x^{p^2} + \cdots + a_nx^{p^n} \in F[x].$$

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### Example

Let  $\mathbb{F}_{125} = \mathbb{F}_5[\theta]/(\theta^3 + \theta + 1)$ .

$$f = x^{25} + (3\theta^2 + 4\theta + 2)x^5 + (3\theta^2 + 4\theta + 2)x$$

is an additive polynomial, and

$$\begin{aligned} f &= (x^5 + (\theta^2 + \theta + 4)x) \circ (x^5 + 3\theta x) \\ &= (x^5 + (2\theta^2 + 4\theta + 2)x) \circ (x^5 + (\theta^2 + 2\theta)x) \end{aligned}$$



## Ore's Legacy

In 1932-4, Oystein Ore wrote four seminal papers for finite fields, differential algebra, and computer algebra

- 1 O. Ore, *Formale Theorie der linearen Differentialgleichungen*, J. reine angew. Math., v. 168, pp. 233-252, 1932.
- 2 O. Ore, *Theory of Non-Commutative Polynomials*, "Annals of Mathematics", v. 34, no. 22, pp. 480-508, 1933.
- 3 O. Ore, *On a Special Class of Polynomials*, Trans. Amer. Math. Soc., v. 35, pp. 559-584, 1933.
- 4 O. Ore, *Contributions to the Theory of Finite Fields*, Trans. Amer. Math. Soc., v. 36, pp. 243-274, 1934.

[1,2] form the basis for modern computational theory of LODEs  
(Ore\_algebra,OreTools)

[3,4] have had great influence on theory of finite fields

Additive polynomials are employed in

- Error correcting codes
- HFE and other cryptosystems
- Mathematical constructions in algebraic function fields
- General fun and parlour tricks.

**Despite their large (exponential) degrees we will see that we can compute very efficiently with them.**

## The Geometry of Additive Polynomials

Denote the set of all additive polynomials over  $\mathbb{F}_q$  as

$$\mathbb{F}_q[x; p] = \left\{ a_0x + a_1x^p + \dots + a_nx^{p^n} \in \mathbb{F}_q[x] \right\}$$

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- Roots  $V_f$  of  $f$  form an  $\mathbb{F}_p$ -vector space of  $\overline{\mathbb{F}}_q$  of dimension  $n$ .
- If  $W$  is any  $\mathbb{F}_p$ -subspace of  $V_f$ , and  $h \in \overline{\mathbb{F}}_q[x]$  has roots exactly  $W$  (i.e.,  $h(W) = 0$ )
  - ➔  $h \in \overline{\mathbb{F}}_q[x; p]$  and  $\exists g \in \overline{\mathbb{F}}_q[x; p]$  such that  $f = g \circ h$ .

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  - ➔  $h \in \overline{\mathbb{F}}_q[x; p]$  and  $\exists g \in \overline{\mathbb{F}}_q[x; p]$  such that  $f = g \circ h$ .
- If  $W$  is also  $\sigma_q$ -invariant, then  $h \in \mathbb{F}_q[x; p]$   
 *$\sigma_q$  is known as the Frobenius automorphism*

## The Geometry of Additive Polynomials (2)

### Example

Again let  $\mathbb{F}_{125} = \mathbb{F}_5[\theta]/(\theta^3 + \theta + 1)$ , and

$$f = x^{25} + (3\theta^2 + 4\theta + 2)x^5 + (3\theta^2 + 4\theta + 2)x$$

Then

$$\mu = \text{RootOf}(x^4 + (\theta^2 + 3\theta + 4)x^2 + (3\theta^2 + 4\theta)x + (4\theta^2 + \theta))$$

$$\nu = \text{RootOf}(x^4 + (4\theta^2 + 2\theta + 1)x^2 + (4\theta^2 + 2\theta)x + (4\theta^2 + \theta))$$

$$V_f = \{\alpha\mu + \beta\nu : \alpha, \beta \in \mathbb{F}_p\} \subseteq \mathbb{F}_{5^{12}}$$

$$\sigma_q = \begin{pmatrix} 3 & 3 \\ 2 & 3 \end{pmatrix} \quad (\text{after some ugly calculations})$$

Probably not the best way to work with additive polynomials...



## Right Composition Factors as EigenVectors of $\sigma_q$

Given  $f \in \mathbb{F}_q[x; p]$  of degree  $n$ , let's find

$$\#\left\{h = x^p + ax \in \mathbb{F}_q[x; p] : \exists g \in \mathbb{F}_q[x; p] \text{ with } f = g \circ h\right\}$$

The number of right composition factors of  $f$  degree  $p$

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The number of right composition factors of  $f$  degree  $p$

= number of 1-dimensional  $\sigma_q$ -invariant subspaces of  $V_f$

= number of eigenvectors of  $\sigma_q$

Remember,  $\sigma_q : V_f \rightarrow V_f$  is a  $\mathbb{F}_p$ -linear map

➡  $\sigma_q$  acts like an  $n \times n$  matrix over  $\mathbb{F}_p$

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New questions:

- How many eigenvectors can an  $n \times n$  matrix over  $\mathbb{F}_q$  have?
- How can we compute this?

## Right Composition Factors as EigenVectors of $\sigma_q$ (2)

How many eigenvectors can a matrix have?

Look at the (rational) Jordan form in  $\mathbb{F}_p^{n \times n}$

**Example:** degree  $p^2$  ( $n = 2$ ): the number of ways of decomposing

$$\begin{aligned} f &= x^{p^2} + a_1 x^p + a_0 x \\ &= (x^p + b_0 x) \circ (x^p + c_0 x) \end{aligned}$$

Put  $\sigma_q$  in rational Jordan form; there are only four possibilities:

$$\sigma_q \sim \begin{pmatrix} 0 & \alpha \\ 1 & \beta \end{pmatrix}, \begin{pmatrix} \lambda & 1 \\ 0 & \lambda \end{pmatrix}, \begin{pmatrix} \lambda & 0 \\ 0 & \mu \end{pmatrix}, \begin{pmatrix} \lambda & 0 \\ 0 & \lambda \end{pmatrix},$$

Here  $\lambda, \mu, \alpha, \beta \in \mathbb{F}_p^*$ ,  $\lambda \neq \mu$  and  $y^2 - \beta y - \alpha \in \mathbb{F}_p[y]$  is irreducible.

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0

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0                      1

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Here  $\lambda, \mu, \alpha, \beta \in \mathbb{F}_p^*$ ,  $\lambda \neq \mu$  and  $y^2 - \beta y - \alpha \in \mathbb{F}_p[y]$  is irreducible.

**An  $f \in \mathbb{F}_q[x; \sigma]$  of degree  $p^2$  can have only 0, 1, 2, or  $p + 1$  right composition factors of degree  $p$ .**



## Right Composition Factors as EigenVectors of $\sigma_q$ (3)

**Example:** degree  $p^3$  ( $n = 3$ ): the number of ways of decomposing

$$f = x^{p^3} + a_2 x^{p^2} + a_1 x^p + a_0 x$$

$$= (x^{p^2} + b_1 x^p + b_0 x) \circ (x^p + c_0 x)$$

$$\sigma_q \sim \begin{pmatrix} \lambda & & \\ & \lambda & \\ & & \lambda \end{pmatrix}, \begin{pmatrix} \lambda & 1 & \\ & \lambda & \\ & & \lambda \end{pmatrix}, \begin{pmatrix} \lambda & 1 & \\ & \lambda & 1 \\ & & \lambda \end{pmatrix}, \begin{pmatrix} \lambda & 1 & \\ & \lambda & \\ & & \mu \end{pmatrix}$$

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$p + 2$ 
 $3$ 
 $1$ 
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$$\sigma_q \sim \begin{array}{cccc} \begin{pmatrix} \lambda & & \\ & \lambda & \\ & & \lambda \end{pmatrix}, & \begin{pmatrix} \lambda & 1 & \\ & \lambda & \\ & & \lambda \end{pmatrix}, & \begin{pmatrix} \lambda & 1 & \\ & \lambda & 1 \\ & & \lambda \end{pmatrix}, & \begin{pmatrix} \lambda & 1 & \\ & \lambda & \\ & & \mu \end{pmatrix} \\ p^2 + p + 1 & p + 1 & 1 & 2 \\ \begin{pmatrix} \lambda & & \\ & \lambda & \\ & & \mu \end{pmatrix}, & \begin{pmatrix} \lambda & & \\ & \mu & \\ & & \nu \end{pmatrix}, & \begin{pmatrix} \lambda & & \\ & \square & \\ & & \square \end{pmatrix}, & \begin{pmatrix} \square & & \\ & \square & \\ & & \square \end{pmatrix} \\ p + 2 & 3 & 1 & 0 \end{array}$$

➔ An  $f \in \mathbb{F}_q[x; \sigma]$  of degree  $p^3$  can have only

$$0, 1, 2, 3, p + 1, p + 2, \text{ or } p^2 + p + 1$$

right composition factors of degree  $p$ .

## General categorization of number of composition factors

How many composition factors of degree  $p$  can an additive polynomial of degree  $p^n$  have?  $S_n$  is the set of possible numbers:

$$S_0 = \{0\}$$

$$S_1 = \{0, 1\}$$

$$S_2 = \{0, 1, 2, p + 1\}$$

$$S_3 = \{0, 1, 2, 3, p + 1, p + 2, p^2 + p + 1\}$$

$$S_4 = \{0, 1, 2, 3, 4, 2p + 2, p^2 + p + 2, p^3 + p^2 + p + 1\}$$

$$\vdots \quad \vdots$$

In general  $\#S_n = \sum_{0 \leq k \leq n} P(k)$ , where  $P(k)$  is the number of additive partitions of  $k$ .

## Efficient Counting of Composition Factors

Roots of  $f \in \mathbb{F}_q[x; p]$  of degree  $p^n$  may be in an extension field of high degree ( $O(p^{O(n^2)})$ ).

➔ Can't really compute directly with  $V_f$ .

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$\mathbb{F}_q[x; p]$  is a (non-commutative) ring under the  $+$  and  $\circ$

- Left (and right) Euclidean ring: LCLM and GCRD operations.
- No unique factorization (but Jordan-Hölder and Krüll Schmidt give a lot of structure to factorizations)
- Fast algorithms for  $+$ ,  $\circ$ , lclm and gcd (time  $O(n^3 \log^2 q)$ ).



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Example:  $\mathbb{F}_{125}[x; 5]$  again – a left Euclidean ring

$$f = x^{25} + (3\theta^2 + 4\theta + 2)x^5 + (3\theta^2 + 4\theta + 2)x$$

$$g = x^{25} + (3\theta^2 + \theta + 3)x^5 + (4\theta^2 + 2\theta + 2)x$$

$$f + g = 2x^{25} + (3\theta^2 + 2\theta + 3)x^5 + (4\theta^2 + 3\theta + 2)x$$

$$f \circ g = x^{625} + (4\theta^2 + 2)x^{125} + \dots + (2\theta^2 + 3\theta + 1)x$$

$$\text{lclm}(f, g) = x^{125} + (\theta^2 + 3\theta + 1)x^{25} + (2\theta^2 + 3)x^5 + (2\theta^2 + 2\theta + 3)x$$

$$\text{gcd}(f, g) = x^5 + 3\theta x$$

## The Centre of Things

The **centre** of  $\mathbb{F}_q[x; p]$  is also very useful:

$$\text{centre}(\mathbb{F}_q[x; p]) = \mathbb{F}_p[x; q] = \left\{ \sum \alpha_i x^{q^i} \in \mathbb{F}_p[x] \right\}$$

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### A cool trick

Given any  $f \in \mathbb{F}_q[x; p]$  we can find a left multiple in the center.

Can do this with  $O(n^3 \log^2 q)$  operations in  $\mathbb{F}_q$ .

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For example (again in  $\mathbb{F}_{125}$ ):

$$\begin{aligned} f &= x^{25} + (3\theta^2 + 4\theta + 2)x^5 + (3\theta^2 + 4\theta + 2)x \in \mathbb{F}_q[x; 5] \\ f^* &= x^{125^2} + 4x^{125} + 3x \in \mathbb{F}_p[x; 125] \end{aligned}$$

*$f^*$  is the minimal central left multiple (mclm) of  $f$*

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The mclm can be found with  $O(n^3 \log^2 q)$  operations in  $\mathbb{F}_q$

## The centre of things (2)

Basis of the factoring algorithm in Giesbrecht (1992, 1998):

Factor the minimal central left multiple and take GCRDs:

$$f = x^{25} + (3\theta^2 + 4\theta + 2)x^5 + (3\theta^2 + 4\theta + 2)x \in \mathbb{F}_q[x; 5]$$

$$f^* = x^{125^2} + 4x^{125} + 3x \in \mathbb{F}_p[x; 125]$$

$$\mapsto y^2 + 4y + 3 = (y + 1)(y + 3)$$

$$f^* = \underbrace{(x^{125} + x)}_{f_1} \circ \underbrace{(x^{125} + 3x)}_{f_2} = (x^{125} + 3x) \circ (x^{125} + x)$$

$$\left. \begin{aligned} \text{gcd}(f, f_1) &= x^5 + (\theta^2 + 2\theta)x \\ \text{gcd}(f, f_2) &= x^5 + 3\theta x \end{aligned} \right\} \text{right composition factors of } f$$

## Factorization in $\mathbb{F}_q[x; p]$

Theorem: (Giesbrecht 1992, 1998)

*Given  $f = \sum_{0 \leq i \leq n} a_i x^{p^i} \in \mathbb{F}_q[x]$ , we can find  $g, h \in \mathbb{F}_q[x]$ , if they exist, such that  $f = g \circ h$ . Requires expected time  $O(n^4 \log^2 q)$  operations in  $\mathbb{F}_q$  (Las Vegas).*



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Hardest when minimal central left multiple is irreducible in  $\mathbb{F}_p[y]$ .

- Construct a finite algebra  $\mathcal{A}$  from  $f$ , called the *eigenring*; show that zero-divisors in  $\mathcal{A}$  yields composition factors of  $f$ .
- Show how to find zero divisors in a finite algebra quickly.
- Build very explicit Krüll-Schmidt and Jordan-Hölder like decompositions, which show structure of all decompositions

## Central Multiples and Frobenius Automorphisms

Theorem: (von zur Gathen, Giesbrecht, and Ziegler 2010)

Let  $f \in \mathbb{F}_q[x; p]$  be squarefree of degree  $p^n$  with roots  $V_f$ , and let  $\sigma_q : V_f \rightarrow V_f$  be the Frobenius automorphism.

Let  $f^* \in \mathbb{F}_p[x; q]$  be the minimal central left multiple of  $f$ .

$f^* = \sum_{0 \leq i \leq m} \alpha_i x^{q^i} \implies f^+ = \sum_{0 \leq i \leq m} \alpha_i y^i$  is min poly of  $\sigma_q$ .

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- ➔ We can find the minimal polynomial of  $\sigma_q$  quickly
- ➔ With a little more work we can compute the complete rational Jordan form of  $\sigma_q$ 
  - ➔ We can count the number of eigenvectors/right factors of degree  $p$  quickly:
  - ➔ Given  $f \in \mathbb{F}_q[x; p]$  of degree  $p^n$ , we can compute the number of right composition factors of degree  $p$  with  $O(n^3 \log^2 q)$  operations in  $\mathbb{F}_q$ .

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Back to our example in  $\mathbb{F}_{125}[x; 5]$

$$f = x^{25} + (3\theta^2 + 4\theta + 2)x^5 + (3\theta^2 + 4\theta + 2)x \in \mathbb{F}_q[x; 5]$$

$$\begin{aligned} f^* &= x^{125^2} + 4x^{125} + 3x \in \mathbb{F}_p[x; 125] \\ &= (x^{125} - 4x) \circ (x^{125} - 2x) \end{aligned}$$

So  $\sigma_q \sim \begin{pmatrix} 4 & 0 \\ 0 & 2 \end{pmatrix}$  and  $\begin{cases} \sigma_q \text{ has two eigenvectors} \\ f \text{ has two right factors of degree 5} \\ h_1 = x^5 + \theta^2 x + 2\theta x, h_2 = x^5 + 3\theta x \end{cases}$

## Projective Polynomials

Projective polynomials were defined by Abhyankar (1997) as a way of constructing polynomials with specific Galois groups.

$$\Psi_n^{(a,b)} = x^{(p^n-1)/(p-1)} + ax + b \in \mathbb{F}_q[x] \text{ for } b \neq 0$$

They have recently been shown to have numerous applications: strong Davenport pairs, difference sets, cryptographically secure sequences, construction of error-correcting codes...

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### Lemma

$\Psi_n^{(a,b)}$  has a root  $c \in \mathbb{F}_q \iff x^{p^n} + ax^p + bx = g \circ (x^p - cx)$ .

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### Theorem

*We can compute the number of roots in  $\mathbb{F}_q$  of  $\Psi_n^{(a,b)} \in \mathbb{F}_q[x]$  with  $O(n^3 \log^2 q)$  operations in  $\mathbb{F}_q$  (even though it has degree  $\approx p^{n-1}$ ).*



## Inverse Problems

How many additive polynomials of degree  $n$  have each possible number of right factors?

Equivalently: how many projective polynomials have each possible number of roots?

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Bluher (2004): For  $f = x^{p^2} + a_1x^p + a_0x \in \mathbb{F}_q[x; p]$

Right factors of degree $p$	# additive polynomials of degree $p^2$ with that many right factors
0	$\frac{p(q^2-1)}{2(p+1)}$
1	$\frac{q^2-q}{p} + 1$
2	$\frac{(q-1)^2 \cdot (p-2)}{2(p-1)} + q - 1$
$p + 1$	$\frac{(q-1)(q-p)}{p(p^2-1)}$

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We give an elementary proof and a way to efficiently enumerate all classes

## Inverse Problems (2)

We now have a general method to give formulas for the number of additive polynomials with a prescribed number of right factors of degree  $p$ .

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Von zur Gathen & Giesbrecht (2011): for

$$f = x^{p^3} + a_2 x^{p^2} + a_1 x^p + a_0 x \in \mathbb{F}_q[x; p]$$

Right factors of degree $p$	# additive polynomials of degree $p^3$ with that many right factors
0	$\frac{1}{3} \frac{(p^3-p)(q^3-1)}{p^3-1}$
1	?
2	?
3	$\frac{(p-2)(p-3)(q-1)^3}{(p-1)^2}$
$p + 1$	?
$p + 2$	$\frac{(q-1)^2 (q-p)(p-2)}{(p^2-1)(p^2-p)}$
$p^2 + p + 1$	$\frac{(q-1)(q-p)(q-p^2)(p-1)}{(p^3-1)(p^3-p)(p^3-p^2)}$

## Conjecture

Consider *any* polynomial  $f \in \mathbb{F}_q[x]$  of degree  $p^2$ .

**Conjecture:**  $f$  can have either  $0, 1, 2, p + 1$  decompositions.

Verified in Sage for  $p \leq 11$ .

In fact, we think they all fall into very specific families.

## Open Questions

- Inverse theory for number of right factors of degree  $p$  of any polynomial in  $\mathbb{F}_q[x; p]$
- Automatically generate inverse formulas
- Compute number of right factors of any given degree of a polynomial in  $\mathbb{F}_q[x; \sigma]$
- Resolve conjecture: how many decompositions possible for a general polynomial?