# Faster Sparse Interpolation of Straight-Line Programs

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# Outline

An overview of straight-line programs
- Straight-line programs
- Statement of the problem
- Summary of results

A new, recursive sparse interpolation algorithm
- "ok" primes - primes that give us information about *some* of the terms of $f$
- Building an approximation with possible errors
- Recursively interpolating the error

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# Straight-line programs

A straight-line program (SLP) is a model of algebraic computation. An SLP can represent a polynomial (a **functional representation**).

## Definition

*Let $R$ be a ring, and let $b_0 \in R$ be a user-specified input. A* **straight-line program (SLP)** *is a series instructions $(\Gamma_1, \ldots, \Gamma_\ell)$, where $\Gamma_i$ assigns a value to $b_i \in R$, and $\Gamma_i$ is of the form*

$$\Gamma_i : b_i \leftarrow \alpha \ [+ - \times] \ \beta, \qquad \alpha, \beta \in R \cup \{b_0, \ldots, b_{i-1}\}.$$

*The vector $b$ is the output corresponding to input $b_0$. We say $\ell$ is the* **length** *of our program.*

## Note

This definition is a restriction to **single-input**, **division-free** SLPs.

# Straight-line programs

## Example: a SLP for $f(z) = (z^3 + z^2 + z + 1)(z - 1)$

$b_0 = z$ is our input.

$$b_1 \leftarrow b_0 \times b_0 \quad [= z^2], \qquad b_2 \leftarrow b_0 \times b_1 \quad [= z^3],$$
$$b_3 \leftarrow b_2 + b_1 \quad [= z^3 + z^2], \qquad b_4 \leftarrow b_3 + b_0 \quad [= z^3 + z^2 + z],$$
$$b_4 \leftarrow b_3 + 1 \quad [= z^3 + z^2 + z + 1], \qquad b_5 \leftarrow b_0 - 1 \quad [= z - 1],$$
$$b_6 \leftarrow b_4 \times b_5 \quad [= (z^3 + z^2 + z + 1)(z - 1)].$$

We say this SLP **computes** $f$. Interpolating this SLP means expanding

$$(z^3 + z^2 + z + 1)(z - 1) = z^4 - 1.$$

# Probing a straight-line program

## The interpolation problem

We are given:

1. A SLP that computes $f(z) = \sum_{i=1}^{t} c_i z^{e_i} \in R[z]$, where
   - $e_1 < e_2 < \cdots < e_t = \deg(f)$, and
   - $c_i \neq 0$ for $1 \leq i \leq s$.
2. Upper bounds $D \geq \deg(f)$ and $T \geq t$.

**Aim**: Construct the terms $c_i z^{e_i}$ of $f$.

- One can **probe** the SLP, i.e., execute the program on chosen inputs.
- Input: $\zeta$, a symbolic $n$-th root of 1, for select $n$.
  - This gives $f(\zeta) \bmod (\zeta^n - 1)$.
  - $n =$ "the **probe degree**".

## Note

In the black-box polynomial model, one instead inputs *all* $n$ $n$-th roots of 1 in order to construct $f(\zeta) \bmod (\zeta^n - 1)$.

# Probing a straight-line program

## Example: Computing $f(\zeta)$ mod $(\zeta^3 - 1)$

$b_0 \leftarrow \zeta$,

$b_1 \leftarrow b_0 \times b_0 \quad [= \zeta^2]$,  $b_2 \leftarrow b_0 \times b_1 \quad [= \zeta^3 = 1]$,

$b_3 \leftarrow b_2 + b_1 \quad [= 1 + \zeta^2]$,  $b_4 \leftarrow b_3 + b_0 \quad [= 1 + \zeta + \zeta^2]$,

$b_4 \leftarrow b_3 + 1 \quad [= 2 + \zeta + \zeta^2]$,  $b_5 \leftarrow b_0 - 1 \quad [= \zeta - 1]$,

$b_6 \leftarrow b_4 \times b_5 \quad [= (2 + \zeta + \zeta^2)(\zeta - 1) = \zeta - 1]$.

Each SLP instruction entails adding/multiplying two polynomials modulo $(\zeta^n - 1)$

**Cost of one SLP instruction**: $\widetilde{\mathcal{O}}(n)$

**Cost of probe**: $\widetilde{\mathcal{O}}(\ell n)$

# Probing a straight-line program

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## Measure of cost of interpolation

cost $\approx$ (# of probes)$*$(probe degree)

# Results of Monte Carlo interpolation algorithms

| algorithm | # of probes | probe degree | "cost" |
|---|---|---|---|
| Garg and Schost[1] | $\widetilde{\mathcal{O}}(T \log D)$ | $\mathcal{O}(T^2 \log D)$ | $\widetilde{\mathcal{O}}(T^3 \log^2 D)$ |
| Giesbrecht and Roche[2] | $\widetilde{\mathcal{O}}(\log D)$ | $\mathcal{O}(T^2 \log D)$ | $\widetilde{\mathcal{O}}(T^2 \log^2 D)$ |
| AA, Giesbrecht, Roche | $\widetilde{\mathcal{O}}(\log T \log D)$ | $\mathcal{O}(T \log^2 D)$ | $\widetilde{\mathcal{O}}(T \log^3 D)$ |

These algorithms admit a failure probability of $\leq \epsilon$, for a fixed parameter $\epsilon$.

## Recall

$f(z)$ is in $\widetilde{\mathcal{O}}(g(z))$ if $f \in \mathcal{O}\left(g(z)\log(g(z))^c\right)$ for some constant $c$.

1. *Interpolation of polynomials given by straight-line programs. Theoretical Comp. Sci. 2009.*
2. *Diversification improves interpolation. ISSAC 2011.*

# From Monte Carlo to Las Vegas

These algorithms are Monte Carlo (probably correct, deterministic run-time). To verify an output $f^*$ we can use the following:

## Lemma (Blaser et al., 2009)

*Let $g = f - f^*$ be a polynomial over an integral domain $R$. Suppose*

1. *We know $g$ has at most $T$ terms and degree at most $D$.*
2. *$g \bmod (z^p - 1) = 0$ for some $(T-1)\log_2 D$ primes $p$.*

*Then $g = 0$.*

- $\#$ of probes: $\mathcal{O}(T \log D)$
- probe degree: $\widetilde{\mathcal{O}}(T \log D)$
- **cost**: $\widetilde{\mathcal{O}}(T^2 \log^2 D)$

## Note

This test is at least as fast as the Garg-Schost and Giesbrecht-Roche algorithms; however, the test is costlier than the new algorithm when $T \in o(\log D)$.

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# A new, recursive sparse interpolation algorithm

Throughout the algorithm, we build a sparse polynomial $f^*$ (initially zero) approximating the $f$ given by our input SLP. In each recursive step of the algorithm we will refine $f^*$, until $f - f^* = 0$.

## Refining our approximation $f^*$

We interpolate the difference $g = f - f^*$ as follows:

1. With high probability, find an "ok" prime, a prime $p$ for which at most a small, specified proportion of the terms of $g$ collide modulo $(z^p - 1)$.
2. Given $p$, look at images of the form $g(z) \bmod (z^{pq} - 1)$, $q \in \mathbb{Z}$, in order to construct a sparse polynomial $f^{**}$ such that $g - f^{**}$ has at most $T/2$ terms.
3. Set $f^* \leftarrow f^* + f^{**}$ and $T \leftarrow \lfloor T/2 \rfloor$ and repeat. If $T = 0$, return $f^*$.

# good primes

- We say two terms $c_1 z^{e_1}$ and $c_2 z^{e_2}$ **collide** modulo $(z^p - 1)$ if $p \mid (e_1 - e_2)$.
- Both the Garg-Schost and Giesbrecht-Roche algorithms require a **good prime** $p$ for which no terms of $f$ collide modulo $(z^p - 1)$.

## Example of a good prime

Let $f(z) = 1 + z + 3z^{10}$. Then

$$f(z) \bmod (z^5 - 1) = 4 + z,$$
$$f(z) \bmod (z^7 - 1) = 1 + z + 3z^3.$$

- 5 is **not** a good prime as 1 collides with $3z^{10}$ modulo $(z^5 - 1)$.
- $f(z) \bmod (z^7 - 1)$ has three terms (like $f$), so 7 is a good prime.

# ok primes

- An ok prime is a weaker notion of a good prime, that allows for a small number of terms to collide.
- Ok primes will allow us to use primes of size $\mathcal{O}(T \log D)$ instead of $\mathcal{O}(T^2 \log D)$.

## Definition

- *Given a polynomial $g$ with $\leq T$ terms, and a prime $p$, we let $\mathcal{C}_g(p)$ denote the number of terms of $f$ that are involved in collisions modulo $(z^p - 1)$.*
- *We call $p$ an* **ok prime** *if $\mathcal{C}_g(p) < \frac{3}{8}T$.*
- *$p$ is a* **good prime** *if $\mathcal{C}_g(p) = 0$.*

# ok primes - an example

Let $g = 1 + z + z^4 - 2z^{13}$.

## $p = 2$

$$\begin{aligned} 1 + z + z^4 - 2z^{13} \bmod (z^2 - 1) &= 1 + z + 1 - 2z, \\ &= 2 - z. \end{aligned}$$

$z^4$ collides with 1 and $-2z^{13}$ collides with $z$, so $\mathcal{C}_g(2) = 4$.

## $p = 3$

$$\begin{aligned} 1 + z + z^4 - 2z^{13} \bmod (z^3 - 1) &= 1 + z + z - 2z, \\ &= 1. \end{aligned}$$

$z$, $z^4$, and $-2z^{13}$ collide, so $\mathcal{C}_g(3) = 3$.

# "ok" primes

## Lemma

*Let $g(z)$ have degree $\leq D$ and $\leq T$ terms, and let*

$$\lambda = \max\left(21, \left\lceil \tfrac{160}{9}(T-1)\ln D \right\rceil\right) \in \mathcal{O}(T \log D).$$

*Then $p$, a prime chosen at random in the range $[\lambda, 2\lambda]$ satisfies $\mathcal{C}_g(p) < \frac{3}{16}\mathbf{T}$ with probability at least $\frac{1}{2}$.*

Thus, if we look at $f(z)$ mod $(z^p - 1)$ for some $\lceil \log 1/\epsilon \rceil$ primes $p \in [\lambda, 2\lambda]$, we will have come across an ok prime $p$ with probablity $\geq 1 - \epsilon$.

## Cost to search for an ok prime

- **Probe degree**: $p \in \mathcal{O}(T \log D)$.
- **# of probes**: $\mathcal{O}(\log 1/\epsilon)$.
- **Cost**: $\mathcal{O}(T \log D \log 1/\epsilon)$.

# "ok" primes

**Problem**

How do we know the $p$ which minimizes $\mathcal{C}_g(p)$?

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**Lemma**

*Suppose* $g \bmod (z^q - 1)$ *has sparsity* $s_q$ *and* $g \bmod (z^p - 1)$ *has sparsity* $s_p \geq s_q$. *Then* $\mathcal{C}_g(p) \leq 2\mathcal{C}_g(q)$.

Thus choosing the $p$ for which the image $g(z) \bmod (z^p - 1)$ has the **most terms** gives us $\mathcal{C}_g(p) < 2\frac{3}{16}T = \frac{3}{8}T$ (with probability $1 - \epsilon$).

# Constructing images of $g = f - f^*$

Given an ok prime $p$, we then construct images $g(z)$ mod $(z^{pq_i} - 1)$ for a set of coprime $q_i > 1$ chosen as follows. Let

$$x = \max(2\ln(D), 17), \text{ and}$$
$$k = \lceil x/\ln(x) \rceil \in \mathcal{O}\left(\frac{\log D}{\log\log D}\right).$$

Then, for $1 \leq i \leq k$, we set

$$q_i \longleftarrow \text{“greatest power of the } i\text{-th prime not exceeding } x\text{”}.$$

Note that $q_i \in \mathcal{O}(\log D)$ and $\prod_{i=1}^{k} q_i > D$.

## Cost of probes used to build $f^{**}$

- **Probe degree**: $pq_i \in \mathcal{O}(T \log^2 D)$.
- **# of probes**: $k \in \widetilde{\mathcal{O}}(\log D)$.
- **Cost**: $\widetilde{\mathcal{O}}(T \log^3 D)$.

# Building $f^{**}$ approximating $g = f - f^{*}$

- If a term $cz^e$ of $g$ avoids collision modulo $(z^p - 1)$, then in any image $g$ mod $(z^{pq} - 1)$ the term $cz^e$ will appear as a unique term of degree congruent to $e$ mod $p$.

- Conversely, suppose there exists $c_0 z^{e_0}$, where $e_0 \leq D$, satisfying
  1. There is a unique term in the reduced image $g$ mod $(z^p - 1)$ that *may* be an image of $c_0 z^{e_0}$.
  2. There is a unique term in each image $g$ mod $(z^{pq_i} - 1)$ that *may* be an image of $c_0 z^{e_0}$,

  ...then $c_0 z^{e_0}$ may be a term of $g$, and we add it to $f^{**}$.

- $c_0 z^{e_0}$ is not always a term of $g$. If not we call $c_0 z^{e_0}$ a **deceptive term.**

# How to build $f^{**}$ approximating $g$

## Example of a deceptive term

Let

$$g(z) = z^2 + z^{40} + z^{60} + \mathbf{z^{11+4}} - \mathbf{z^{14\cdot 11+4}} - \mathbf{z^{15\cdot 11+4}},$$

and let $p = 11$ and consider $q = 2, 3, 5$. We have

$$\begin{aligned}
g(z) \bmod (z^{11} - 1) &= z^2 + z^7 + z^5 - \mathbf{z^4}, \\
g(z) \bmod (z^{22} - 1) &= z^2 + z^{18} + z^{16} - \mathbf{z^{15}}, \\
g(z) \bmod (z^{33} - 1) &= z^2 + z^7 + z^{27} - \mathbf{z^{26}}, \\
g(z) \bmod (z^{55} - 1) &= z^2 + z^{40} + z^5 - \mathbf{z^{48}}.
\end{aligned}$$

- The first three terms of each image correspond to the terms $z^2, z^{40}, z^{60}$ appearing in $g$. Note there is only one term of degree $\{2, 40, 60\}$ mod 11 in each image.
- The remaining term in each of the 4 images has degree congruent to 4 mod 11. By Chinese remaindering on the exponents, this gives a deceptive term $-\mathbf{z^{323}}$ not appearing in $g$.

# Detecting collisions

- A deceptive term can only result from a collision of **three or more** terms modulo $(z^p - 1)$.
- If only two terms collide modulo $z^p - 1$, there will be a $q_i$ such that $g \bmod (z^{pq_i} - 1)$ separates those terms.

**Example:** $g(z) = 1 + z^{59} + z^{11+2} + z^{11\cdot 12+11+2}$

Using $p = 11$ and $(q_1, q_2, q_3) = (4, 3, 5)$, we have

$$g(z) \bmod (z^{11} - 1) = 1 + 2z^2 + z^4$$
$$g(z) \bmod (z^{44} - 1) = 1 + z^{15} + 2z^{13}$$
$$g(z) \bmod (z^{33} - 1) = 1 + z^{26} + 2z^{13}$$
$$g(z) \bmod (z^{55} - 1) = 1 + z^4 + z^{13} + z^{35}$$

We recognize that a collision occurred at degree 2 modulo $(z^{11} - 1)$ and ignore those terms. Here $f^{**}$ would be $1 + z^{59}$.

# Building an approximation $f^{**}$ of $g = f - f^{**}$

- The polynomial $f^{**}$ we construct will contain the $T - \mathcal{C}_g(p)$ non-colliding terms of $g$, plus potentially some $\lfloor \mathcal{C}_g(p)/3 \rfloor$ deceptive terms.
- $g - f^{**}$ will have at most $\mathcal{C}_g(p) + \frac{1}{3}\mathcal{C}_g(p) = \frac{4}{3}\mathcal{C}_g(p)$ terms.
- If $\mathcal{C}_g(p) < \frac{3}{8}T$, then $g - f^{**}$ will have sparsity less than $(\frac{4}{3})(\frac{3}{8})T = T/2$.

# Recursively interpolating $g - f^{**}$

- $f^{**}$ gives us a new difference $f - f^* - f^{**}$ with smaller sparsity bound $T/2$.
- We set $f^* \leftarrow f^* + f^{**}$ and $T \leftarrow \lfloor T/2 \rfloor$, then recursively interpolate our now updated difference $g = f - f^*$.
- We continue in this fashion some $\lfloor 1 + \log T \rfloor$ times until we have $f - f^* = 0$.

**Problem:**

At the start of our algorithm, we look at $\lceil \log 1/\epsilon \rceil$ primes $p_i$ in an attempt to find an ok prime with probability $1 - \epsilon$. We now need this to succeed at each of the $\lfloor 1 + \log T \rfloor$ recursive calls now.

If we want to correctly interpolate $f$ with probability $\mu$, it suffices to instead bound the failure probability at each recursive step by $\epsilon = \mu/(1 + \log T)$. This does not affect the "soft-Oh" cost of the algorithm.

**Cost with probability of success at least $1 - \mu$, for fixed $\mu$**

Interpolating $f$ entails $\widetilde{\mathcal{O}}(\log T \log D)$ probes of degree $\mathcal{O}(T \log^2 D)$.
**Cost**: $\widetilde{\mathcal{O}}(T \log^3 D)$

- Final thoughts:
  - Difficult to guarantee zero collisions for small probe degree - birthday paradox.
  - ok primes: better performance by tolerating *some* errors.
  - An advantage of having a recursive algorithm decrementing $T$ is that we can call the Giesbrecht-Roche $T^2 \log^2 D$ algorithm once $\log D \gg T$.

- Future work:
  - Investigate the numerical stability of a black-box variant of the algorithm.
  - Las Vegas algorithm: faster polynomial identity testing of SLPs.
  - $\mathcal{O}(T \log D)$ interpolation?

Thank you for your attention