# Gaston, Maple and Mike 

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GNOME 2014, Zurich, July 4th, 2014

Me, Gaston and Maple

May 1982 - Dec 1982 Waterloo, Masters student Jan 1983 - Aug 1989 Waterloo, PhD student Aug 1989 - Oct 1995 Zurich, Assistent

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Gaston gave me this paper for my Masters essay
Shafi Goldswasser and Silvio Micali.
Probabilistic encryption \& how to play mental poker keeping secret all partial information. STOC '82, June 1982
which we implemented in Maple.

Gaston's number theory package, the first Maple package.

```
    |\\~/| Maple V Release 4 (WMI Campus Wide License)
._|\| |/I_. Copyright (c) 1981-1996 by Waterloo Maple Inc. All rights
    \ MAPLE / reserved. Maple and Maple V are registered trademarks of
    <_-_- _-_-> Waterloo Maple Inc.
    Type ? for help.
> with(numtheory);
Warning, new definition for order
[ F, M, cyclotomic, divisors, factorset, fermat, ifactor, imagunit,
    isprime, issqrfree, ithprime, jacobi, lambda, legendre, mcombine,
    mersenne, mlog, mroot, msqrt, nextprime, order, phi, prevprime,
    pprimroot, primroot, quadres, rootsunit, safeprime, sigma, tau]
```

I chose not to pursue cryptography for a PhD.

## Life as a graduate student with Gaston ...



First Maple retreat, Sparrow lake, summer, 1983

## What was Gaston's main contribution to Maple?

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Maple's Sum-of-Products representation and hashing of all subexpressions.


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9 x y^{3} z-4 y^{3} z^{2}-6 x y^{2} z-8 x^{3}-5
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What is the most important operation to make efficient?
Polynomial multiplication (and division).
But monomial multiplication cost $>200$ cycles.

Singular's representation


Our new POLY representation (default in Maple 17)


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6 advantages
(1) It's about $4 \times$ more compact.
(2) Memory access is sequential.
(3) Kernel operations become $O$ (\#terms), some $O(1)$.
(9) Monomial multiplication is one 64 bit integer + Monomial comparison is one 64 bit integer >
(3) The simpl table is not filled with PRODs.
(0) Division cannot cause exponent overflow in graded lex order.

What will fast multiplication using POLY do for the Maple library?

Intel Core i7 920 2.66 GHz (4 cores)
Times in seconds

| multiply | Maple | Maple 16 |  | $\begin{array}{r} \hline \text { Magma } \\ 2.16-8 \\ \hline \end{array}$ | $\begin{array}{r} \text { Singular } \\ 3.1 .0 \end{array}$ | Mathem atica 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 13 | 1 core | 4 cores |  |  |  |
| $p_{4}:=f_{4}\left(f_{4}+1\right)$ | 95.97 | 2.14 | 0.643 | 13.25 | 30.64 | 273.01 |
| divide |  |  |  |  |  |  |
| $q_{4}:=p_{4} / f_{4}$ | 192.87 | 2.25 | 0.767 | 18.54 | 14.96 | 228.83 |
| factor | Hensel lifting is mostly polynomial multiplication! |  |  |  |  |  |
| $p_{4} 135751$ terms | 2953.54 | 59.29 | 46.41 | 332.86 | 404.86 | 655.49 |

$$
f_{4}=(1+x+y+z+t)^{20}+1 \quad 10626 \text { terms }
$$

Parallel speedup for $f_{4} \times\left(f_{4}+1\right)$ is $2.14 / .643=3.33 \times$. Why?

What will fast multiplication using POLY do for the Maple library?

Intel Core i7 9202.66 GHz (4 cores)
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Parallel speedup for $f_{4} \times\left(f_{4}+1\right)$ is $2.14 / .643=3.33 \times$. Why? Conversion overhead between POLY and SUM of PRODs!

After brainstorming with Roman, I asked Laurent if we could make POLY the default in Maple. Maple 17 uses POLY if all monomials in a polynomial with integer coefficients fit in 64 bits - otherwise we use SUM-of-PRODs. Conversions between POLY and SUM-of-PRODs are automatic and invisible to the Maple user.


So we coded POLY for each kernel routine.
Faster at everything except op, map, etc.

| command | Maple 16 | Maple 17 | speedup | notes |
| :--- | :--- | :--- | :--- | :--- |
| $\operatorname{coeff}(f, x, 20)$ | 2.140 s | 0.005 s | 420 x | terms easy to locate |
| $\operatorname{coeffs}(f, x)$ | 0.979 s | 0.119 s | 8 x | reorder exponents and radix |
| degree $(f, x)$ | 0.073 s | 0.003 s | 24 x | stop early using monomial d |
| $\operatorname{diff}(f, x)$ | 0.956 s | 0.031 s | 30 x | terms remain sorted |
| eval $(f, x=6)$ | 3.760 s | 0.175 s | 21 x | use Horner form recursively |
| expand $(2 * x * f)$ | 1.190 s | 0.066 s | 18 x | terms remain sorted |
| indets $(f)$ | 0.060 s | 0.000 s | $\rightarrow O(1)$ | first word in dag |
| op $(f)$ | 0.634 s | 2.420 s | 0.26 x | has to construct old structur |
| for t in f do | 0.646 s | 2.460 s | 0.26 x | has to construct old structur |
| taylor $(f, x, 50)$ | 0.668 s | 0.055 s | 12 x | get coefficients in one pass |
| type $(f$, polynom $)$ | 0.029 s | 0.000 s | $\rightarrow O(n)$ | type check variables only |
| $f ;$ | 0.162 s | 0.000 s | $\rightarrow O(n)$ | evaluate the variables |

For $f$ with $n=3$ variables and $t=10^{6}$ terms created by
$\mathrm{f}:=\operatorname{expand}(m u l(r a n d p o l y(v$, degree $=100$, dense $), v=[x, y, z])$ ):

|  Maple 16  Maple 17  <br> multiply 1 core 4 cores 1 core 4 cores <br> $p_{4}:=f_{4}\left(f_{4}+1\right)$ 2.140 0.643 1.770 0.416 <br> factor     <br> $p_{4} 135751$ terms 59.27 46.41 24.35 12.65 |
| :--- |

$$
f_{4}=(1+x+y+z+t)^{20}+1 \quad 10626 \text { terms }
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Parallel speedup for $f_{4} \times\left(f_{4}+1\right)$ is $1.77 / 0.416=4.2 \times$. How ?

Joris van der Hoven: Do you use the extra bits for the total degree? My answer: No, because ...
I changed my mind. Roman Pearce recoded everything for Maple 18.

|  | per variable |  | total degree |  |  | $V_{n}=$ det $n \times n$ Vandermonde |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $n$ | \#bits | maxdeg | \#bits | maxdeg | deg | Maple 16 | 17 | 18 |
| 7 | 8 | 255 | 8 | 255 | 21 | 0.012 s | 0.005 | 0.004 |
| 8 | 7 | 127 | 8 | 255 | 28 | 0.093 s | 0.027 | 0.026 |
| 9 | 6 | 63 | 10 | 1023 | 36 | 1.35 s | 0.218 | 0.150 |
| 10 | 5 | 31 | 14 | 16383 | 45 | 15.95 s | 25.44 | 1.57 |
| 11 | 5 | 31 | 9 | 511 | 55 | - | - | 18.87 |
| 12 | 4 | 15 | 16 | 65535 | 66 |  |  | 236.4 |
| 13 | 4 | 15 | 12 | 4095 | 78 |  |  | - |
| 14 | 4 | 15 | 8 | 255 | 91 |  |  |  |
| 15 | 4 | 15 | 4 | 15 | 105 |  |  |  |
| 16 | 3 | 7 | 16 | 65535 | 120 |  |  |  |



Maple retreat, Sparrow lake, circa 1992
Thank you Gaston for Waterloo, Zurich and Maple. Mike.

## Notes on integration of POLY for Maple 17

Given a polynomial $f\left(x_{1}, x_{2}, \ldots, x_{n}\right)$, we store $f$ using POLY if
(1) $f$ is expanded and has integer coefficients,
(2) $d>1$ and $t>1$ where $d=\operatorname{deg} f$ and $t=$ \#terms,
(3) we can pack all monomials of $f$ into one 64 bit word, i.e. if $d<2^{b}$ where $b=\left\lfloor\frac{64}{n+1}\right\rfloor$
Otherwise we use the sum-of-products representation.

- The representation is invisible to the Maple user. Conversions are automatic.
- POLY polynomials will be displayed in sorted order.
- Packing is fixed by $n=\#$ variables.
- Maple 18 uses remaining bits for total degree.


## Parallel multiplication using a binary heap.



Target architecture

Local Heaps


One thread per core.

Threads write to a finite circular buffer.


Threads try to acquire global heap as buffer fills up to balance load.

