

Lec12 Solving First order Recurrences

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Lecture 12 Solving First Order Recurrence Relations

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Grimaldi Chapter 10.1

Midterm average 68.6
median 71.3

$$p_n = n p_{n-1}$$

$$c_n = (n-1) + c_{n-1}$$

$$h_n = 1 + 2h_{n-1}$$

Assignment 3 posted.

Due Monday Feb 22nd @ 11pm.

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1: void Bubblesort( double A[], int n ) {
2: // sort the array A of size n into ascending order
3:   int i; double t;
4:   if( n==1 ) return; ← C1=0.
5:   for( i=1; i<=n-1; i++ )
6:     if( A[i-1] > A[i] ) { line 6 is executed n-1 times.
7:       t = A[i-1]; A[i-1] = A[i]; A[i] = t;
8:     }
9:   Bubblesort(A,n-1); recursive call to Bubblesort to sort the
10:  return; first n-1 entries of A. This does
11: } Cn-1 comparisons.

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$i++ \quad \bar{c} = \bar{c} + 1.$

What should we count to determine the cost of the Bubblesort algorithm?
We will count the number of comparisons between elements of A in line 6.
Let c_n be the number of comparisons.

$$\begin{aligned}
n=1 & \quad c_1 = 0 & \leftarrow \text{line 6} & \leftarrow \text{line 9.} \\
n>1 & \quad c_n = n-1 + c_{n-1}
\end{aligned}$$

This is the same RR as k_n the # of edges in K_n .
 $k_n = k_{n-1} + n-1$ and $k_1 = 0.$

A First Look at Solving Recurrence Relations.

How can we solve RRs like

(1) $b_n = 2b_{n-1}$ for $n \geq 2$ and $b_1 = 2$.

(1) $c_n = c_{n-1} + (n-1)$ for $n \geq 2$ and $c_1 = 0$.

$b_n = 2^n$
 $b_n = \#$ of binary strings of length n .

Example $b_n = 2b_{n-1}$

$$\begin{aligned}
 n \rightarrow n-1 &\Rightarrow b_n = 2^1 b_{n-1} && 1+n-1 = n \\
 &= 2(2b_{n-2}) = 2^2 b_{n-2} && 2+n-2 = n \\
 &= 2^2(2 \cdot b_{n-3}) = 2^3 b_{n-3} && 3+n-3 = n \\
 &= 2^3(2 \cdot b_{n-4}) = 2^4 b_{n-4} && 4+n-4 = n \\
 &\vdots && \\
 &= 2^{n-1} \cdot b_1 = 2^{n-1} \cdot 2 = 2^n && \boxed{}+1 = n.
 \end{aligned}$$

We can check that our solution $b_n = 2^n$ satisfies the RR by substituting it into the RR and IV:

$$\begin{array}{lcl}
 b_1 = 2 & b_n = 2b_{n-1} & b_1 = 2 \\
 b_n = 2^n & 2^n = 2 \cdot 2^{n-1} & \checkmark \\
 b_{n-1} = 2^{n-1} & &
 \end{array}$$

Example $c_n = c_{n-1} + (n-1)$ $c_1 = 0$.

$$\begin{aligned}
 n=2 & \begin{aligned} c_n &= c_{n-1} + n-1 \\ c_{n-1} &= c_{n-2} + n-2 \\ c_{n-2} &= c_{n-3} + n-3 \\ &\vdots \end{aligned} \\
 & \begin{aligned} c_2 &= c_1 + 1 \\ c_1 &= 0 \end{aligned}
 \end{aligned}$$

Add these equations.

$$\begin{aligned}
 \sum_{i=1}^n i &= \frac{n(n+1)}{2} \\
 \sum_{i=1}^{n-1} i &= \frac{(n-1)n}{2}
 \end{aligned}$$

$$\begin{aligned}
 c_n &= n-1 + n-2 + n-3 + \dots + 1 = \sum_{i=1}^{n-1} i = \frac{n(n-1)}{2}
 \end{aligned}$$

Check $c_n = c_{n-1} + n-1$

$$c_n = \frac{n(n-1)}{2}; \quad \frac{n(n-1)}{2} = \frac{(n-1)(n-2)}{2} + (n-1) = (n-1) \left[\frac{n-2}{2} + 1 \right] = (n-1) \cdot \frac{n}{2}$$

Theorem (A different way to solve first order RRs)

Every sequence x_0, x_1, x_2, \dots satisfying the recurrence $x_n = dx_{n-1}$ has the general solution $x_n = cd^n$ for some constant c . (The sequence is a geometric progression.)

$$b_n = 2 \cdot b_{n-1}$$

Proof (substitution)

$$x_n = d \cdot x_{n-1}$$

$$\text{Sol: } x_n = c \cdot d^n : c \cdot d^n = d \cdot (c \cdot d^{n-1})$$

$$\Rightarrow x_{n-1} = c \cdot d^{n-1}$$

This suggests the following general strategy for solving RRs:

- (1) Find the **general solution** to the RR. This will have one or more constants.
Note: a RR of order k will has k constants.
- (2) Use the k initial values to determine the constants. This gives a **unique** solution.

Example. Suppose $x_n = 5x_{n-1}$ and $x_0 = 7$. First find the general solution then the unique solution satisfying $x_0 = 7$.

$$n=0 : \quad \begin{aligned} x_n &= d \cdot 5^n \quad \text{for some constant } d. \\ x_0 &= d \cdot 5^0 = d. \\ x_0 &= 7. \end{aligned} \Rightarrow d = 7.$$

$$x_n = 7 \cdot 5^n$$

Exercise 1. Solve $p_n = np_{n-1}$ where $p_1 = 1$.

$$\begin{aligned} p_n &= n \cdot p_{n-1} = n(n-1)p_{n-2} = n(n-1)(n-2) \cdot p_{n-3} \\ &= \dots \\ p_{n-1} &= (n-1) \cdot p_{n-2} \\ p_{n-2} &= (n-2) \cdot p_{n-3} \\ &= n(n-1) \cdot \dots \cdot 3 \cdot 2 \cdot p_1 \\ &= n(n-1) \cdot \dots \cdot 3 \cdot 2 \cdot 1 \\ &= n! \end{aligned}$$

Exercise 2. Solve $x_n = x_{n-1} + An + B$ for $x_1 = C$.

Q2(a). How many strings of length 4 over $\{A, B, C, D\}$ have BB in them?

any of A, B, C, D

$$\begin{array}{ccc} \underline{B} & \underline{B} & \downarrow \downarrow \\ & & 4 \cdot 4 = 16 \end{array} \quad \begin{array}{ccc} \downarrow & \underline{B} & \underline{B} & \downarrow \\ & & & 4 \cdot 4 = 16 \end{array} \quad \begin{array}{ccc} \downarrow & \downarrow & \underline{B} & \underline{B} \\ & & & 4 \cdot 4 = 16 \end{array}$$

total = $3 \cdot 16 = 48$.
This is incorrect.

any of A, C, D

$$\begin{array}{ccc} \underline{B} & \underline{B} & \downarrow \downarrow \\ & & 3 \cdot 3 = 9 \end{array} \quad \begin{array}{ccc} \downarrow & \underline{B} & \underline{B} & \downarrow \\ & & & 3 \cdot 3 = 9 \end{array} \quad \begin{array}{ccc} \downarrow & \downarrow & \underline{B} & \underline{B} \\ & & & 3 \cdot 3 = 9 \end{array}$$

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A, C, D

$$\begin{array}{ccc} \underline{B} & \underline{B} & \underline{B} & \downarrow \\ & & & 3 \end{array} + \begin{array}{ccc} \underline{B} & \downarrow & \underline{B} & \underline{B} \\ & & & 3 \end{array} + \begin{array}{ccc} \downarrow & \underline{B} & \underline{B} & \underline{B} \\ & & & 3 \end{array} + \begin{array}{ccc} \underline{B} & \underline{B} & & \downarrow \\ & & & 3 \end{array} + \begin{array}{ccc} \underline{B} & & & \downarrow \\ & & & 3 \end{array}$$

$$\begin{array}{ccc} \underline{B} & \underline{B} & \underline{B} & \underline{B} \\ & & & 1 \end{array}$$

$$\begin{array}{r} 27 \\ 3 = 12 \\ \hline = 1 \\ \hline 40 \end{array}$$

Q2 (b) How many ways can three teams of 5 players be selected from 15 players?

$$\rightarrow \binom{15}{5} \cdot \binom{10}{5} \cdot \binom{5}{5} / 6 = 3! \text{ ways.}$$

team 1 team 2 team 3

one way 1,3,5,7,9 2,4,6,8,10 11,12,13,14,15

same teams. 2,4,6,8,10 1,3,5,7,9 11,12,...,15

There are $3!$ ways to permute the 3 teams so we need to divide by $3! = 6$.