

# Lecture 29: Trees and Rooted Trees

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Grimaldi 12.1, 12.2

## Definition ( tree )

A multigraph  $G$  is a **tree** if  $G$  is connected and  $G$  does not contain a cycle.

## Theorem ( main properties of trees )

*If  $T = (V, E)$  is a tree then  $|V| = |E| + 1$  and secondly, there is a unique path in  $T$  between every pair of vertices.*

Examples

## Theorem (Characterization of Trees)

*Let  $G = (V, E)$  be a multigraph. The following statements are equivalent.*

- (1)  $G$  is connected and has no cycle. ( $G$  is a tree)*
- (2)  $G$  is connected and  $|V| = |E| + 1$ .*
- (3)  $G$  has no cycle and  $|V| = |E| + 1$ .*
- (4) There is a unique path between every pair of vertices in  $G$ .*

Proof.

Proof (cont).

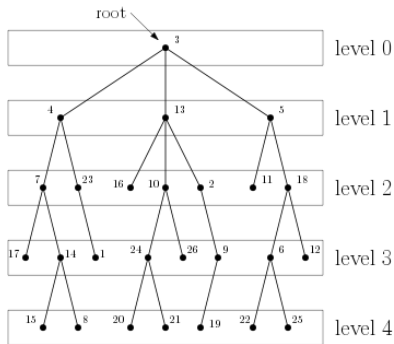
## Definition ( rooted tree )

A **rooted tree**  $T = (V, E)$  is a tree with a distinguished vertex called the **root**. For every vertex  $v \in V$  the **level** of  $v$  is the length of the path from  $v$  to the root. Note: the root is the unique vertex at level 0.

Example

## Definition ( rooted tree terminology )

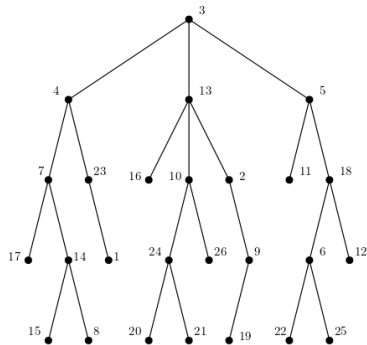
- The **height** of a rooted tree is the maximum level of a vertex. A rooted tree consisting of just a root vertex has height 0.
- Every non-root vertex  $v$  at level  $i$  is adjacent to exactly one vertex  $u$  at level  $i - 1$ . We call  $u$  the **parent** of  $v$  and we say that  $v$  is a **child** of  $u$ .
- For every vertex  $v$  there is a walk “up the tree” to the root obtained by moving to the parent vertex at each step. If  $u$  is another vertex on this walk, we call  $u$  an **ancestor** of  $v$  and  $v$  a **descendant** of  $u$ .



We are frequently interested in working with rooted trees recursively. Therefore, it will be helpful to think of a rooted tree as composed out of smaller rooted trees.

## Definition ( subtree )

Let  $v$  be a vertex of a rooted tree  $T$  with level  $i$ . Define  $T'$  to be the subgraph of  $T$  induced by  $v$  together with its descendants. Then  $T'$  forms a new rooted tree with root vertex  $v$ . We say that  $T'$  with root  $v$  is the **subtree** of  $T$  at  $v$ .



## Definition ( isomorphism of rooted trees)

Let  $T_1, T_2$  be rooted trees with  $T_i = (V_i, E_i)$  for  $i = 1, 2$ . We say that  $T_1$  and  $T_2$  are **isomorphic** if there exists a bijection  $f : V_1 \rightarrow V_2$  satisfying:

- (1)  $\{f(u), f(v)\} \in E_2 \Leftrightarrow \{u, v\} \in E_1$
- (2) For every  $v \in V_1$  the level of  $v$  and  $f(v)$  is the same.  
In particular,  $f$  sends the root of  $T_1$  to the root of  $T_2$ .

Example.



## Definition

A rooted tree is  **$m$ -ary** if every internal node has at most  $m$  children. A 2-ary tree is called **binary** tree.

Exercise. Find all binary trees with height 0, 1, and 2 up to isomorphism.

Let  $b_n$  denote the number of binary trees of height at most  $n$ . Find  $b_0$ ,  $b_1$ ,  $b_2$ .

Use the recursive structure of rooted trees find a recurrence for  $b_n$ .