

Lecture 27 Trees

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Grimaldi 12.1

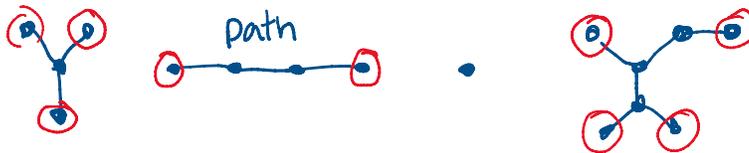
Assignment #7 is due Wednesday.
Upto today's lecture.

Definition (tree, forest and leaf)

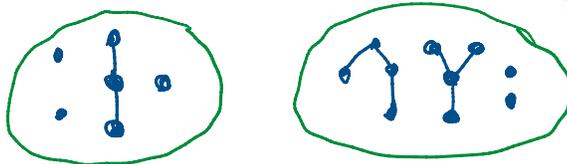
Let $G = (V, E)$ be a multigraph. G is a **tree** if G is connected and G does not contain a cycle. G is a **forest** if G does not contain a cycle. A vertex of degree 1 is called **leaf** or **pendant vertex**.

Examples

Trees



Forests



Since a tree cannot have loops or parallel edges, it is a simple graph.

We previously showed that every graph with all vertices of degree ≥ 2 must have a cycle. Therefore every tree with ≥ 2 vertices must contain a leaf. Later we will see

that trees have ≥ 2 leaves.

A tree is a forest.
A forest may not be a tree.

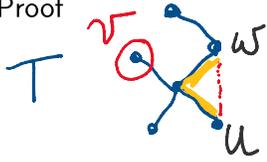
Lemma

If $T = (V, E)$ is a tree with leaf v then $T - v$ is a tree.

Show

[$T - v$ is connected and acyclic]

Proof



First observe if $T - v$ has a cycle. then T has a cycle which would contradict T is a tree.

Let $u, w \in V$ with $u \neq w \neq v$. There is a path from u to w in T (as T is connected). The only two vertices of degree 1 in a path are the end vertices.

All other vertices have degree ≥ 2 .

So v (has degree 1) is not on the path.

Therefore in $T - v$ there is a path from u to w .

Therefore $T - v$ is connected.

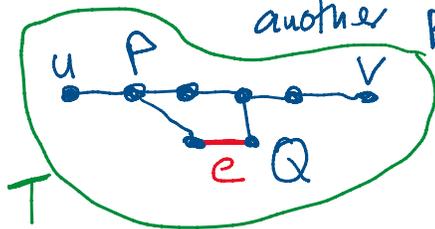
This observation gives us a powerful tool for proving properties of trees. Try using induction on the number of vertices and, for the inductive step, deleting a leaf then applying the inductive hypothesis.

Theorem (unique paths)

If $T = (V, E)$ is a tree and $u, v \in V$ are distinct, there is a unique path in T with ends u, v .

Proof. (is a path from u to v) A tree is connected so there is a path from u to v .

(path is unique). Towards a contradiction suppose there is another path Q from u to v . Since $Q \neq P$ there must



be at least one edge $e \in E$ that is on Q but not on P . Let $e = \{a, b\}$.

Notice $T - e$ is connected. Therefore e is in a cycle in T . This contradicts T is a tree.

So there is a unique path from u to v .

Theorem (main property of trees)

If $T = (V, E)$ is a tree then $|V| = |E| + 1 \Rightarrow |E| = |V| - 1$.

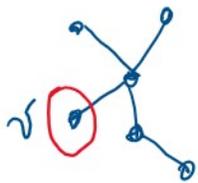
If $G = (V, E)$ is a forest with k trees then $|V| = |E| + k$.

Proof. By induction on $|V|$. Let $n = |V|$ in T .

Base: $n=1$ T is the singleton vertex $\bullet \stackrel{=1}{=} \overset{=0}{=} \checkmark$
 Here $|V|=1$, $|E|=0$ and $|V|=|E|+1 \checkmark$

Ind. Step: $n \geq 2$:

Ind. Hypothesis. Assume $|V|=|E|+1$ holds for any tree with $< n$ vertices.



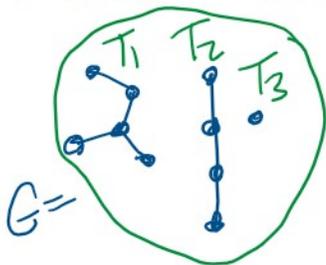
Let v be a leaf in T .
 Then $T-v$ is a tree with $|V|-1$ vertices and $|E|-1$ edges. By induction

$$T-v \text{ satisfies } (|V|-1) = (|E|-1) + 1 \Rightarrow |V| = |E| + 1.$$

Proof (cont).

By induction on n , $|V|=|E|+1$ holds for all trees with n vertices.

Forrests. Let T_1, T_2, \dots, T_k be the trees in the forest G .



By the first part of the theorem

$$|V_i| = |E_i| + 1 \text{ for } 1 \leq i \leq k \text{ where } T_i = (V_i, E_i)$$

Summing over i we get

$$\underbrace{|V_1| + |V_2| + \dots + |V_k|}_{= |V|} = \underbrace{(|E_1| + 1) + (|E_2| + 1) + \dots + (|E_k| + 1)}_{|E| + k}$$

So $|V| = |E| + k$ in a forest of k trees.

Lemma

If $G = (V, E)$ satisfies $|V| = |E| + 1$ then G must have a vertex of degree 0 or at least two of degree 1. $\Rightarrow |E| = |V| - 1$.



Proof. Let $n = |V|$ and let k_0 be the # vertices of degree 0 and k_1 be the # of vertices of degree 1 in G . We want to show that $k_0 \geq 1$ or $k_1 \geq 2$ (or both). Consider the degree sum formula from 11-4

$$2|E| = \sum_{v \in V} \deg(v)$$

$$2|E| = 2(|V| - 1) \Rightarrow 0 \cdot k_0 + 1 \cdot k_1 + 2(\underbrace{n - k_0 - k_1}_{\text{\# vertices of degree } \geq 2})$$

$$2(n-1) \geq k_1 + 2n - 2k_0 - 2k_1$$

$$\cancel{2n} - 2 \geq \cancel{2n} + k_1 - 2k_0 - \cancel{2k_1}$$

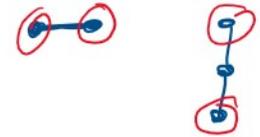
$$k_1 + 2k_0 \geq 2. \Rightarrow k_1 \geq 2 \text{ or } k_0 \geq 1.$$

at least degree 2.

vertices of degree ≥ 2

Lemma

Every tree $T = (V, E)$ with $|V| \geq 2$ has at least two leaves.



Proof. T is a tree $\Rightarrow T$ is connected
 and $T \neq \bullet \Rightarrow$ every vertex has degree ≥ 1 .
 \Rightarrow there are no vertices of degree 0
 \Rightarrow there are ≥ 2 vertices of degree 1 (previous Lemma)
 \Rightarrow there are two leaf vertices.

Summary. If T is a tree with $n = |V|$ vertices

- If $n \geq 2$ then T has ≥ 2 leaves
- If T has 2 leaves then T is a path
- T has $n - 1$ edges.

$$|V| = |E| + 1$$

$$|E| = |V| - 1$$

$$|E| = n - 1$$