

[Assignment 3 Question 1 (on the FFT)

▼ Part (a)

```
> p := 97;
                                     p:= 97
> ifactor(p-1);
                                     (2)5 (3)
> primedivisors := [2,3];
                                     primedivisors:= [2, 3]
> alpha := 5;
                                     α:= 5
> for q in primedivisors do modp( alpha((p-1)/q), p ); od;
                                     96
                                     35
Neither power is 1 so α = 5 is a primitive element.
> w := alpha((p-1)/8) mod p;
                                     w:= -33
> `mod` := mods;
                                     mod:= mods
> seq( wi mod p, i=0..8 );
                                     1, -33, 22, -47, -1, 33, -22, 47, 1
```

▼ Part (b)

What is the Fourier transform of the polynomial $a(x) = 1 + x + x^2 + \dots + x^{n-1}$. Well let's just try it

```
> a := add( xi, i=0..7 );
                                     a:= x7 + x6 + x5 + x4 + x3 + x2 + x + 1
> F := [seq( eval(a,x=wi) mod p, i=0..7 )];
                                     F:= [8, 0, 0, 0, 0, 0, 0, 0]
```

So the answer must be $[n, 0, 0, 0, \dots, 0]$. Obviously, for $\omega = 1$ we have $a(1) = n$. Recall that for $x \neq 1$, we have

$1 + x + x^2 + \dots + x^{n-1} = \frac{(1 - x^n)}{1 - x}$ which Maple seems to "know".

```
> simplify( sum(xi,i=0..n-1)*(1-x) );
                                     -xn + 1
```

So that the Fourier transform entry $F_i = \frac{(1 - \omega^{in})}{(1 - \omega^i)}$ but the numerator is 0.

Part (c)

$$M(n) = 2 \cdot M\left(\frac{n}{2}\right) + \frac{n}{2} \text{ hence}$$

$$2 \cdot M\left(\frac{n}{2}\right) = 4 \cdot M\left(\frac{n}{4}\right) + \frac{2 \cdot n}{4} = 4 \cdot M\left(\frac{n}{4}\right) + \frac{n}{2} \text{ hence}$$

$$4 \cdot M\left(\frac{n}{4}\right) = 4 \cdot M\left(\frac{n}{8}\right) + \frac{n}{2} \text{ hence ... hence}$$

$$2^{k-1} \cdot M\left(\frac{n}{2^k}\right) = 2^{k-1} \cdot M\left(\frac{n}{2^k}\right) + \frac{n}{2} \text{ and finally}$$

$$2^k \cdot M\left(\frac{n}{2^k}\right) = 2^k \cdot M(1) = 0.$$

Adding both sides and cancelling equal terms we have $k = \log_2(n)$ lots of $\frac{n}{2}$ on the right and $M(n)$ on the left that don't cancel. Hence $M(n) = \frac{n}{2} \cdot \log_2(n)$.

Part (d)

Katatsuba recurrence

```
> simplify( rsolve( {T(1)=d, T(n)=3*T(n/2)+c*n}, T(n) ) );
```

$$(2c + d) n^{\frac{\ln(3)}{\ln(2)}} - 2cn$$

Optimized FFT recurrence (see code below)

```
> simplify( rsolve( {T(1)=0, T(n)=2*T(n/2)+n/2}, T(n) ) );
```

$$\frac{1}{2} \frac{n \ln(n)}{\ln(2)}$$

Gaussian elimination recurrence

```
> simplify( rsolve( {T(1)=0, T(n)=T(n-1)+(n-1)^2}, T(n) ) );
```

$$\frac{1}{6} n + \frac{1}{3} n^3 - \frac{1}{2} n^2$$

Part (e)

```
> unprotect(FFT):
```

```
FFT := proc(n,A,p,w)
```

```
local n2,B,C,i,wi,T;
```

```
if n=1 then return; fi;
```

```
n2 := n/2;
```

```
B := Array(0..n2-1);
```

```
C := Array(0..n2-1);
```

```
for i from 0 to n2-1 do B[i] := A[2*i]; C[i] := A[2*i+1]; od;
```

```
FFT(n2,B,p,w^2 mod p);
```

```
FFT(n2,C,p,w^2 mod p);
```

```
wi := 1;
```

```
for i from 0 to n2-1 do
```


Part (f)

Now let's multiply $a(x) \cdot b(x)$ using the FFT.

```
> a := 1+3*x-x^3;
  b := 2*x^4-3*x^3-2*x^2+x+1;
      a := -x^3 + 3x + 1
      b := 2x^4 - 3x^3 - 2x^2 + x + 1
<
> A := Array(0..7,[1,3,0,-1]) mod p;
> B := Array(0..7,[1,1,-2,-3,2]):
> W := Array(0..3,[seq(w^i mod p,i=0..3)]):
> FFT(8,A,p,W,1);
> FFT(8,B,p,W,1);
> C := Array(0..7):
> for i from 0 to 7 do C[i] := A[i]*B[i] mod p: od:
> W := Array(0..3,[seq(1/w^i mod p,i=0..3)]):
> FFT(8,C,p,W,1);
> ninv := 1/8 mod p:
> for i from 0 to 7 do C[i] := ninv*C[i] mod p od:
> add( C[i]*x^i, i=0..7 );
      -2x^7 + 3x^6 + 8x^5 - 8x^4 - 10x^3 + x^2 + 4x + 1
<
Let's check
> expand(a*b) mod p;
      -2x^7 + 3x^6 + 8x^5 - 8x^4 - 10x^3 + x^2 + 4x + 1
```