

Sparse Polynomials

November 23, 2023 7:06 PM

Let $f \in R[x_1, \dots, x_n]$, R a ring.

Let $f = \sum_{i=1}^t a_i \cdot M_i(x_1, \dots, x_n)$ where $a_i \in R$, $a_i \neq 0$,
 \uparrow
 monomials.

If $\deg(f, x_i) = d_i$ then f may have upto $t = \prod_{i=1}^n (d_i + 1)$ terms.

E.g. $n=2$, $d_1=3$, $d_2=2$. $t \leq (3+1)(2+1) = 12$.

		deg x_2		
		0	1	2
deg x_1	0	1	x_2	x_2^2
	1	x_1	$x_1 x_2$	$x_1 x_2^2$
	2	x_1^2	$x_1^2 x_2$	$x_1^2 x_2^2$
	3	x_1^3	$x_1^3 x_2$	$x_1^3 x_2^2$

We say f is sparse if
 $t = \#f \ll \prod_{i=1}^n (d_i + 1)$.

? $t \leq \sqrt{\prod_{i=1}^n (d_i + 1)}$.

E.g. $f = 3x_1 + 2x_1 x_2 + 5x_2 x_1^2$

Alternatively if $\deg(f) = d$ then f may have upto $\binom{n+d}{d}$ terms.

E.g. $d=3$, $n=2$ $\binom{n+d}{d} = \binom{2+3}{3} = \frac{5!}{2!3!} = \frac{5 \cdot 4}{2!} = 10$ terms.

1	x_1	x_1^2	x_1^3
	x_2	$x_1 x_2$	$x_1^2 x_2$
	x_2^2	$x_1 x_2^2$	
	x_2^3		

We say f is sparse if
 $t \ll \binom{n+d}{d}$.

I use $t \leq \sqrt{\binom{n+d}{d}} = \sqrt{10}$

So $f = 1 - x^3 - y^3$ is sparse.

$$T_3 = \begin{bmatrix} u & v & w \\ v & u & v \\ w & v & u \end{bmatrix}$$

$$\det(T_3) = u^3 - 2uv^2 - uw^2 + 2v^2w. \quad t = \underline{4}$$

$$\binom{n+d}{d} = \binom{6}{3} = 20. \quad \uparrow \text{sparse.}$$

In still most polynomials with $n \geq 3$ are sparse.

Sparse Polynomial Interpolation.

Let $f \in k[x_1, \dots, x_n]$ where k is a field.

Suppose f is given by a black-box $B: k^n \rightarrow k$.

Suppose $f = \sum_{i=1}^t a_i M_i(x_1, \dots, x_n)$.

How many points in k^n do we need to interpolate f ?
I.e. to recover the $a_i \in k$ and the monomials M_i .

Let $d_i = \deg(f, x_i)$.

Since $\#f \leq \prod_{i=1}^n (d_i + 1) = D$ we may need D values of f .

E.g. $f(x_1, x_2) = \sum_{j=0}^{d_1} \left(\sum_{i=0}^{d_2} a_{ij} x_2^i \right) \cdot x_1^j$

Pick $\beta_0, \beta_1, \dots, \beta_{d_2} \in k$.

Interpolate $f(x_1, \beta_0) = \bullet + \bullet x_1 + \dots + \bullet x_1^{d_1}$ using $d_1 + 1$ points for x_1
 $f(x_1, \beta_1) = \bullet + \bullet x_1 + \dots + \bullet x_1^{d_1}$ using $d_1 + 1$ points for x_1
 \vdots
 $f(x_1, \beta_{d_2}) = \bullet + \bullet x_1 + \dots + \bullet x_1^{d_1}$ using $d_1 + 1$ points.

$\Rightarrow f(x_1, x_2) = f_0(x_2) + f_1(x_2) \cdot x_1 + \dots + f_{d_1}(x_2) \cdot x_1^{d_1}$ So $(d_2 + 1)(d_1 + 1)$ points

Suppose f is sparse. Can we improve on this?

Dense method. $\# \text{values of } f = \prod_{i=1}^n (d_i + 1) = (1+d)^n \leftarrow \text{is exponential in } n$
 $f = 1 \cdot x_1^d + 1 \cdot x_2^d + \dots + 1 \cdot x_n^d$

Zippel ^{PhD} 1979 $\leq (t \sum_{i=1}^n d_i) + 1 = tnd + 1$

Ben-Or/Tiwari ^{PhD} 1988 $2T$ where $T \geq t$ $2T$