## MACM 401/MATH 701/MATH 819 Assignment 4, Spring 2007.

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This assignment is to be handed in by Tuesday March 13th before the start of class. For problems involving Maple calculations and Maple programming, you should submit a printout of a Maple worksheet of your Maple session.
Late Penalty: $-20 \%$ for each day late.

## Question 1: Resultants (10 marks)

(a) Let $A, B$ be non-zero polynomials in $\mathbb{Z}[x]$. Let $\operatorname{deg} A=m$ and $\operatorname{deg} B=n$ and let $c$ be an integer. Let $\operatorname{res}(A, B)$ denote the resulant of $A$ and $B$. Complete the identities.

$$
\begin{gathered}
\operatorname{res}(A, B)=? ? \operatorname{res}(B, A) \text { and } \\
\operatorname{res}(c A, B)=? ? \operatorname{res}(A, B)
\end{gathered}
$$

(b) Let $A, B$ be two non-zero polynomials in $\mathbb{Z}[x]$. Let $A=G \bar{A}$ and $B=G \bar{B}$ where $G=\operatorname{gcd}(A, B)$. Recall that a prime $p$ in the modular gcd algorithm is unlucky if $p \mid R$ where $R=\operatorname{res}(\bar{A}, \bar{B})=0$ is the resultant of $\bar{A}$ and $\bar{B}$, an integer.
Consider the following pair of polynomials (from assignment 4)

$$
\begin{gathered}
A=58 x^{4}-415 x^{3}-111 x+213, \quad \text { and } \\
B=69 x^{3}-112 x^{2}+413 x+113
\end{gathered}
$$

They are relatively prime, i.e., $G=1, \bar{A}=A$ and $\bar{B}=B$. Using Maple compute the resultant $R$ and identify all unlucky primes. For each unlucky prime $p$ compute the gcd of the polynomials $A$ and $B$ modulo $p$ to verify that the primes are indeed unlucky.

## Question 2: $P$-adic Lifting (20 marks)

## Reference: Section 6.3.

(a) By hand, determine the $p$-adic representation of the integer $u=116$ for $p=5$ using the positive representation, then the symmetric representation for $\mathbb{Z}_{p}$.

Using Maple, determine the $p$-adic representation for the polynomial

$$
u(x)=28 x^{2}+24 x+58
$$

with $p=3$ using, first the positive representation for $\mathbf{Z}_{3}$, then the symmetric representation.
(b) Determine the cube-root, if it exists, of the following polynomials

$$
\begin{aligned}
a(x) & =x^{6}-531 x^{5}+94137 x^{4}-5598333 x^{3}+4706850 x^{2}-1327500 x+125000 \\
b(x) & =x^{6}-406 x^{5}+94262 x^{4}-5598208 x^{3}+4706975 x^{2}-1327375 x+125125
\end{aligned}
$$

using reduction mod 5 and linear $p$-adic lifting. Factor the polynomials so you know what the answers are. Express the answer in the p-adic representation. To calculate the initial solution $u_{0}=\sqrt[3]{a} \bmod 5$ use any method. Use Maple to do the calculations.

## Question 3: Hensel lifting (20 marks)

Reference: Section 6.4 and 6.5.
(a) Given

$$
a(x)=x^{4}-2 x^{3}-233 x^{2}-214 x+85
$$

and image polynomials

$$
u_{0}(x)=x^{2}-3 x-2 \quad \text { and } \quad w_{0}(x)=x^{2}+x+3
$$

satisfying $a \equiv u_{0} w_{0}(\bmod 7)$, lift the image polynomials using Hensel lifting to find (if there exist) $u$ and $w$ in $\mathbb{Z}[x]$ such that $a=u w$.
(b) Given

$$
b(x)=48 x^{4}-22 x^{3}+47 x^{2}+144
$$

and an image polynomials

$$
u_{0}(x)=x^{2}+4 x+2 \quad \text { and } \quad w_{0}=x^{2}+4 x+5
$$

satisfying $b \equiv 6 u_{0} w_{0}(\bmod 7)$, lift the image polynomials using Hensel lifting to find (if there exist) $u$ and $w$ in $\mathbb{Z}[x]$ such that $b=u w$.

## Question 4: Determinants (20 marks)

Consider the following 3 by 3 matrix $A$ of polynomials in $\mathbb{Z}[x]$ and its determinant $d$.

```
> P := () -> randpoly(x,degree=2,dense):
> A := linalg[randmatrix] (3,3,entries=P);
```

$$
A:=\left[\begin{array}{ccc}
-55-7 x^{2}+22 x & -56-94 x^{2}+87 x & 97-62 x \\
-83-73 x^{2}-4 x & -82-10 x^{2}+62 x & 71+80 x^{2}-44 x \\
-10-17 x^{2}-75 x & 42-7 x^{2}-40 x & 75-50 x^{2}+23 x
\end{array}\right]
$$

> d := linalg[det] (A);

$$
d:=-224262-455486 x^{2}+55203 x-539985 x^{4}+937816 x^{3}+463520 x^{6}-75964 x^{5}
$$

(a) Let $A$ by an $n$ by $n$ matrix of polynomials in $\mathbb{Z}[x]$ and let $d=\operatorname{det}(A)$. Develop a modular algorithm for computing $d=\operatorname{det}(A) \in \mathbb{Z}[x]$. Your algorithm will compute determinants of $A$ modulo a sequence of primes and apply the CRT. For each prime $p$ it will compute the determinant in $\mathbb{Z}_{p}[x]$ by evaluation and interpolation. In this way we reduce computation of a determinant of a matrix over $\mathbb{Z}[x]$ to many computations of determinants of matrices over $\mathbb{Z}_{p}$, a field, for which ordinary Gaussian elimination, which costs $O\left(n^{3}\right)$ arithmetic operations in $\mathbb{Z}_{p}$, may be used.
You will need to bounds for $\operatorname{deg} d$ and $\|d\|_{\infty}$. Use primes $p=[101,103,107, \ldots]$ and use Maple to do Chinese remaindering. Use $x=1,2,3, \ldots$ for the evaluation points and use Maple for interpolation. Implement your algorithm in Maple and test it on the above example.

To reduce the coefficients of the polynomials in $A$ modulo $p=7$ in Maple use

```
> p := 7: phip := proc(f) f mod p end:
> B := map( phip, A );
```

To evaluate the polynomials in $B$ at $x=a l p h a(\bmod p)$ in Maple use

```
> alpha := 1: phix := proc(f) Eval(f,x=alpha) mod p end:
```

> C := map( phix, B );

To compute the determinant of a matrix $C$ over $\mathbb{Z}_{p}$ in Maple use

```
> Det( C ) mod p;
```

(b) Suppose $A$ is an $n$ by $n$ matrix over $\mathbb{Z}[x]$ and $A_{i, j}=\sum_{k=0}^{d} a_{i, j, k} x^{k}$ and $\left|a_{i, j, k}\right|<B^{m}$. That is $A$ is an $n$ by $n$ matrix of polynomials of degree at most $d$ with coefficients at most $m$ base $B$ digits long. Assume the primes satisfy $B<p<2 B$ and that arithmetic in $\mathbb{Z}_{p}$ costs $O(1)$. Estimate the time complexity of your algorithm in big $O$ notation as a function of $n, m$ and $d$. Make reasonable simplifying assumptions such as $n<B$ and $d<B$ as necessary. Also helpful is $\ln n!<n \ln n$. State your assumptions.

