

Assignment #5 is due on Monday — end of MACM 401.

Resultant Calculus

Let $A, B \in K[x]$, K a field e.g. \mathbb{Q}

Let $A = a_m x^m + \dots + a_0 = a_m \prod_{i=1}^m (x - \alpha_i)$ where $\alpha_i \in \bar{K}$.

$B = b_n x^n + \dots + b_0 = b_n \prod_{i=1}^n (x - \beta_i)$ where $\beta_i \in \bar{K}$.

$$\text{res}(A, B) = \underline{a_m} \cdot b_n \overset{m}{\prod_{i=1}^m} \overset{n}{\prod_{j=1}^n} (\alpha_i - \beta_j) = (-1)^{mn} b_n a_m \overset{n}{\prod_{j=1}^n} \overset{m}{\prod_{i=1}^m} (\beta_j - \alpha_i)$$

① $\text{res}(c \in K, B) = c^n \cdot b_n^0 = c^n$

② $\text{res}(B, A) = (-1)^{mn} \text{res}(A, B)$.

$$\text{res}(A, B) = a_m^n \prod_{i=1}^m \left(b_n \prod_{j=1}^n (\alpha_i - \beta_j) \right) = (-1)^{mn} b_n^m \prod_{j=1}^n \left(a_m \prod_{i=1}^m (\beta_j - \alpha_i) \right)$$

$B(\alpha_i)$ $A(\beta_j)$

③ $\text{res}(A, B) = a_m^n \prod_{i=1}^m B(\alpha_i) = (-1)^{mn} b_n^m \prod_{j=1}^n A(\beta_j)$.

$\text{deg } m \uparrow$ $\text{deg } n \downarrow$

④ $\text{res}(A, BC) = a_m^{n+e} \prod_{i=1}^m (BC)(\alpha_i) = a_m^n \prod_{i=1}^m B(\alpha_i) \cdot a_m^e \prod_{i=1}^m C(\alpha_i)$

$\text{deg } e \uparrow$

③ $= \text{res}(A, B) \cdot \text{res}(A, C)$.

⑤ $\text{res}(A, B) = 0 \iff \text{deg}(\text{gcd}(A, B)) > 0$.

Suppose $\begin{matrix} A & B \\ m \geq n > 0 \end{matrix}$. Let $R = \text{rem}(A \div B) \implies \boxed{A = BQ + R}$

CASE $R=0 \implies \text{deg}(\text{gcd}(A, B)) > 0 \implies \text{res}(A, B) = 0$.

CASE $R \neq 0$.

Let $l = \deg R$.

$$\text{res}(A, B) = \text{res}(BQ + R, B)$$

$$\textcircled{3} = (-1)^{mn} b_n^m \prod_{j=1}^n \underbrace{(BQ)(\beta_j)}_{=0} + R(\beta_j)$$

$$\Rightarrow \text{res}(A, B) = (-1)^{mn} b_n^m \prod_{j=1}^n R(\beta_j)$$

$$\text{res}(R, B) \textcircled{3} = (-1)^{ln} b_n^l \prod_{j=1}^n R(\beta_j)$$

$$\frac{\text{res}(A, B)}{\text{res}(R, B)} = (-1)^{mn - ln} \cdot b_n^{m-l}$$

$$\textcircled{6} \quad \text{res}(A, B) = (-1)^{n(m-l)} \cdot b_n^{m-l} \text{res}(R, B).$$
$$\text{gcd}(A, B) = \text{gcd}(R, B)$$

We can compute the resultant using the E.A. in $O(n \cdot m)$ arithmetic operations which is better than using $\det(S)$ where S is Sylvester's matrix because that costs $O((n+m)^3)$ using Gaussian elimination.